

Quantum Hypergraph Homomorphisms and Operator Algebras

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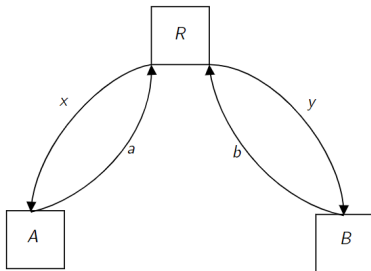
Motivation

Functional analytic methods \rightsquigarrow study combinatorial structures

- **Non-local game**: discrete object we study using **operator algebraic techniques**
- Similarity between non-local games? Strategy transport? Allow comparison for chance of “winning”?
- **Quantum homomorphisms of discrete structures** \rightsquigarrow studied for graphs, but few others
 - Can we do the same for non-local games?

Non-local games

- The (classical) definition of a **non-local game** is a tuple (X, Y, A, B, λ) where X, Y, A, B are finite sets and $\lambda : X \times Y \times A \times B \rightarrow \{0, 1\}$ is a “verifier” function which encodes the rules of the game.



Hypergraphs

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- Reformulate non-local games: a non-local game on (V_2, W_1, V_1, W_2) is a hypergraph $\Lambda \subseteq V_2 W_1 \times V_1 W_2$.
 - So Λ corresponds to the **support of λ** in classical definition.

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- A channel $\mathcal{E} : \mathcal{D}_V \rightarrow \mathcal{D}_W$ defines a hypergraph

$$E_{\mathcal{E}}(w) = \{(v, w) \in V \times W : \mathcal{E}(w|v) > 0\}.$$

(Here $\mathcal{E}(w|v) = \langle \mathcal{E}(\epsilon_{v,v}), \epsilon_{w,w} \rangle$ where $\langle \cdot, \cdot \rangle$ is the **trace** of the matrix product and $\epsilon_{v,v}$ are the basis elements for \mathcal{D}_V .)

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- Channel \mathcal{E} , hypergraph $E_{\mathcal{E}} \iff$ channel \mathcal{E}^* , hypergraph $E_{\mathcal{E}^*} = (E_{\mathcal{E}})^*$

Correlations

Let V_i, W_i be finite sets with $i = 1, 2$. A **no-signalling (NS) correlation** on the quadruple (V_2, W_1, V_1, W_2) is an information channel $\Gamma : \mathcal{D}_{V_2 W_1} \rightarrow \mathcal{D}_{V_1 W_2}$ for which **marginal channels**

$$\Gamma_{V_2 \rightarrow V_1} : \mathcal{D}_{V_2} \rightarrow \mathcal{D}_{V_1}, \quad \Gamma_{V_2 \rightarrow V_1}(v_1 | v_2) := \sum_{w_2 \in W_2} \Gamma(v_1, w_2 | v_2, w_1'),$$

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The collection of no-signalling correlations is denoted by \mathcal{C}_{ns} ; other classes of correlations $(\mathcal{C}_{\text{loc}}, \mathcal{C}_{\text{q}}, \mathcal{C}_{\text{qa}}, \mathcal{C}_{\text{qc}})$ are defined by additional restrictions we place on $\Gamma \in \mathcal{C}_{\text{ns}}$.

Local correlations: $\Gamma \in \mathcal{C}_{\text{loc}}$ is a convex combination of correlations

$$\Gamma(v_1 w_2 | v_2 w_1) = \Gamma_1(v_1 | v_2) \Gamma_2(w_2 | w_1),$$

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Quantum commuting: $\Gamma \in \mathcal{C}_{\text{qc}}$ if

$$\Gamma(v_1 w_2 | v_2 w_1) = \langle E_{v_2 v_1} F_{w_1 w_2} \xi, \xi \rangle$$

for mutually commuting **POVM**'s $(E_{v_2 v_1})_{v_1 \in V_1}, (F_{w_1 w_2})_{w_2 \in W_2}$ acting on \mathcal{H} and $\xi \in \mathcal{H}$ is a unit vector.

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$$\mathcal{C}_{\text{loc}} \subset \mathcal{C}_{\text{q}} \subset \mathcal{C}_{\text{qa}} \subset \mathcal{C}_{\text{qc}} \subset \mathcal{C}_{\text{ns}}$$

Simulation paradigm

For an NS correlation Γ on (V_2, W_1, V_1, W_2) and a channel $\mathcal{E} : \mathcal{D}_{V_1} \rightarrow \mathcal{D}_{W_1}$, the map $\Gamma[\mathcal{E}] : \mathcal{D}_{V_2} \rightarrow \mathcal{D}_{W_2}$ defined by

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- We “wire” the **output for the marginal channel $\Gamma_{V_2 \rightarrow V_1}$** to the input for \mathcal{E} , and the **output of \mathcal{E} back into Γ** .

$$\begin{array}{ccc} V_1 & \xrightarrow{\mathcal{E}} & W_1 \\ \uparrow & & \downarrow \\ V_2 & \xrightarrow{\Gamma[\mathcal{E}]} & W_2 \end{array}$$

When $\Gamma[\mathcal{E}] \in \mathcal{C}(V_2 \times W_2)$ where Γ is the **simulator**, we write $(V_1 \mapsto W_1) \xrightarrow{\Gamma} (V_2 \mapsto W_2)$.

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- For $t \neq \text{ns}$, $\Gamma \in \mathcal{C}_t^{\text{bi}}$ now has slight additional restriction: POVM's $(E_{v_2, v_1})_{v_1, v_2} \in V$ and $(F_{w_1, w_2})_{w_1, w_2} \in W$ are **magic squares**.

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- **Note:** $\mathcal{C}_t(\Lambda) = \mathcal{C}(\Lambda) \cap \mathcal{C}_t, \mathcal{C}_t^{\text{bi}}(\Lambda) = \mathcal{C}(\Lambda) \cap \mathcal{C}_t^{\text{bi}}$.

Definition

We say that

- E_1 is t -homomorphic to E_2 (denoted $E_1 \rightarrow_t E_2$) if $\mathcal{C}_t(E_1 \leftrightarrow E_2) \neq \emptyset$.
- E_1 is t -isomorphic to E_2 (denoted $E_1 \simeq_t E_2$) if $V_1 = V_2, W_1 = W_2$ with $\mathcal{C}_t^{\text{bi}}(E_1 \leftrightarrow E_2) \neq \emptyset$.

Local homomorphisms

A map $f: V_2 \rightarrow V_1$ is a (classical) **homomorphism between hypergraphs** E_1 and E_2 if pre-images under f preserve edge relations; that is, f is a homomorphism if there exists a map $g: W_1 \rightarrow W_2$ so that

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 - If Γ is perfect for $E_1 \rightarrow E_2$, assume Γ is an extreme point in $\mathcal{C}_{\text{loc}} + \text{no-signalling} \rightsquigarrow$ homomorphism (f, g) .
 - If (f, g) a homomorphism between hypergraphs, let $\Phi: \mathcal{D}_{V_2} \rightarrow \mathcal{D}_{V_1}, \Psi: \mathcal{D}_{W_1} \rightarrow \mathcal{D}_{W_2}$ where $\Phi(v_1|v_2) = \delta_{v_1, f(v_2)}$ and $\Psi(w_2|w_1) = \delta_{w_2, g(w_1)}$. Then $\Gamma = \Phi \otimes \Psi \in \mathcal{C}_{\text{loc}}(E_1 \leftrightarrow E_2)$.

An operator system approach

Start with a finite set V , and a block operator matrix $U = (u_{v,v'})_{v,v' \in V}$ such that U and U^t are **isometries**. Let \mathcal{V}_V be the (universal) ternary ring of operators generated by $u_{v,v'}$ for $v, v' \in V$ and the relations

$$\sum_{a \in V} [u_{a'',x''}, u_{a,x}, u_{a,x'}] = \delta_{x,x'} u_{a'',x''}, \quad \sum_{x \in V} [u_{a'',x''}, u_{a,x}, u_{a',x}] = \delta_{a,a'} u_{a'',x''}.$$

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- For a **faithful ternary representation** $\theta : \mathcal{V}_V \rightarrow \mathcal{B}(\mathcal{H}, \mathcal{K})$ (where \mathcal{H}, \mathcal{K} are Hilbert spaces), for the **right C*-algebra** \mathcal{C}_V we have $\mathcal{C}_V \simeq \overline{\text{span}}(\theta(\mathcal{V}_V)^* \theta(\mathcal{V}_V))$.
- Write $e_{v_1, v'_1, v_2, v'_2} := u_{v_2, v_1}^* u_{v'_2, v'_1}$, $v_1, v_2, v'_1, v'_2 \in V$.
 - The C*-algebra \mathcal{C}_V is generated by elements e_{v_i, v'_i, v_2, v'_2} for $v_i, v'_i \in V, i = 1, 2$.

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- For a **faithful ternary representation** $\theta : \mathcal{V}_V \rightarrow \mathcal{B}(\mathcal{H}, \mathcal{K})$ (where \mathcal{H}, \mathcal{K} are Hilbert spaces), for the **right C^* -algebra** \mathcal{C}_V we have $\mathcal{C}_V \simeq \overline{\text{span}}(\theta(\mathcal{V}_V)^* \theta(\mathcal{V}_V))$.
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Theorem (H.-Todorov, in prep. 2022)

The map $s \mapsto \Gamma_s$ is an affine surjective correspondence between

- *the states of $\mathcal{S}_V \otimes_{\max} \mathcal{S}_W$ which annihilate \mathcal{J} and the perfect ns-strategies of $E_1 \leftrightarrow E_2$.*
- *the states of $\mathcal{S}_V \otimes_c \mathcal{S}_W$ which annihilate \mathcal{J} and the perfect qc-strategies of $E_1 \leftrightarrow E_2$.*
- *the states of $\mathcal{S}_V \otimes_{\min} \mathcal{S}_W$ which annihilate \mathcal{J} and the perfect qa-strategies of $E_1 \leftrightarrow E_2$.*

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$$\Gamma_s(v_1, w_2 | v_2, w_1) = s(e_{v_2, v_1} \otimes f_{w_1, w_2})$$

gives us the correspondence with perfect t -strategies on $E_1 \leftrightarrow E_2$ (for $t \in \{\text{qa}, \text{qc}, \text{ns}\}$).

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Remark: When \mathcal{H} is a Hilbert space, a **quantum magic square** over V on \mathcal{H} is a block operator matrix $(E_{v_2, v_1})_{v_1, v_2 \in V}$ with positive entries, and

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Operator system \mathcal{S}_V is **universal for quantum magic squares**:

ucp maps $\phi : \mathcal{S}_V \rightarrow \mathcal{B}(\mathcal{H}) \leftrightarrow$ quantum magic square $(E_{v_1, v_2})_{v_1, v_2 \in V}$ via $E_{v_1, v_2} = \phi(e_{v_1, v_2})$

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Let $t \in \{\text{loc}, \text{q}, \text{qc}\}$. *TFAE:*

- E_1 is faithfully t -isomorphic to E_2 ;
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Note: The ideas for this proof were adapted from Atserias et. al ([5] 2019), where a similar result was shown for graphs only.

Local vs. quantum strategies

For a given finite graph G with vertex set X , we can form hypergraphs

$$E_G = \{(x, x') : x \sim x'\}, \quad F_G = \{((x, y), y) : x \sim_G y\}$$

in $X \times X$ and $XX \times X$, respectively.

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Local vs. quantum strategies

Proof: (Sketch)

(i) As $G_1 \cong_q G_2$, find permutation $P \in M_X \otimes M_d$ intertwining $A_{G_1} \otimes I_d$ and $A_{G_2} \otimes I_d$; this implies $E_{G_1} \cong_q E_{G_2}$.

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(ii) Using permutation $P = (P_{x,y})_{x,y}$ as before, we know

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If $\xi \in \mathcal{H} \otimes \mathcal{H}$ is maximally entangled, set

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Then p gives us a perfect approximately quantum strategy for $F_{G_1} \cong F_{G_2}$.

Strong no-signalling correlations

We restrict ourselves to considering non-local games as hypergraphs. We assume:

- X_i, Y_i, A_i, B_i are finite sets, $E_i \subseteq X_i Y_i \times A_i B_i$, $i = 1, 2$ are non-local games.
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A channel $\Gamma : \mathcal{D}_{X_2 Y_2 \times A_1 B_1} \rightarrow \mathcal{D}_{X_1 Y_1 \times A_2 B_2}$ is **strongly no-signalling (SNS)** if

$$\sum_{b_2 \in B_2} \Gamma(x_1 y_1, a_2 b_2 | x_2 y_2, a_1 b_1) = \sum_{b_2 \in B_2} \Gamma(x_1 y_1, a_2 b_2 | x_2 y_2, a_1 b_1'), \quad b_1, b_1' \in B_1,$$

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- Operator matrix $P = (P_{xy,ab})$ is **NS** if marginal operators $P_{xa} = \sum_b P_{xy,ab}$ and $P^{yb} = \sum_a P_{xy,ab}$ are well-defined.

SNS correlation classes

A NS operator matrix $P = (P_{xy,ab})_{xy,ab}$ is **dilatable** if there is an isometry $V: \mathcal{H} \rightarrow \mathcal{K}$ between Hilbert spaces and mutually commuting POVM's $(E_{xa})_{a \in A}$, $(F_{yb})_{b \in B}$ on \mathcal{K} with

$$P_{xy,ab} := V^* E_{xa} F_{yb} V, \quad x \in X, y \in Y, a \in A, b \in B.$$

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- **approximately quantum** if Γ is a limit of quantum SNS correlations.
- **local** if Γ is quantum and individual entries in operator matrices P, Q commute with themselves as well.

Homomorphisms of non-local games

Theorem (H.-Todorov, in prep. 2022)

Let Γ be an SNS correlation over the quadruple $(X_2 Y_2, A_1 B_1, X_1 Y_1, A_2 B_2)$ and \mathcal{E} be an NS correlation over (X_1, Y_1, A_1, B_1) . Then

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Note: For $\Gamma \in \mathcal{C}_{\text{sqc}}, \mathcal{E} \in \mathcal{C}_{\text{qc}}$ case, say $\Gamma(x_1 y_1, a_2 b_2 | x_2 y_2, a_1 b_1) = \langle P_{x_2, x_1} P^{y_2, y_1} Q_{a_1, a_2} Q^{b_1, b_2} \xi, \xi \rangle$, and $\mathcal{E}(a_1, b_1 | x_1, y_1) = \langle E_{x_1, a_1} F_{y_1, b_1} \eta, \eta \rangle$ where $\xi \in \mathcal{H}, \eta \in \mathcal{K}$ are unit vectors, and families of operators are mutually commuting POVM's on resp. Hilbert spaces. Set

$$\tilde{E}_{x_2, a_2} = \sum_{x_1 \in X_1} \sum_{a_1 \in A_1} P_{x_2, x_1} Q_{a_1, a_2} \otimes E_{x_1, a_1}, \quad \tilde{F}_{y_2, b_2} = \sum_{y_1 \in Y_1} \sum_{b_1 \in B_1} P^{y_2, y_1} Q^{b_1, b_2} \otimes F_{y_1, b_1}.$$

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We then have **qc-decomposition** $\Gamma[\mathcal{E}](a_2, b_2 | x_2, y_2) = \langle \tilde{E}_{x_2, a_2} \tilde{F}_{y_2, b_2} (\xi \otimes \eta), \xi \otimes \eta \rangle$. (Others follow similarly).

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– Holds for SNS bicorrelations as well.

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- Non-local games are now **t-isomorphic** if we can find perfect SNS strategies $\Gamma \in \mathcal{C}_t^{\text{bi}}(E_1 \leftrightarrow E_2)$.
- If $E_1 \rightarrow_{\text{st}} E_2$ or $E_1 \simeq_{\text{st}} E_2$, we simulate optimal strategies for E_1 using SNS bicorrelation Γ and get strategies for E_2 .

Thank you for listening!

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