



MapReduce and its use for indexing

The Programming Model and Practice

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- MapReduce programming model
 - **Brief intro to MapReduce**
 - Use of MapReduce inside Google
 - MapReduce programming examples
 - MapReduce, similar and alternatives
 - Practical indexing examples in IR
 - Inverted index construction
 - PageRank computation
 - Implementation of Google MapReduce
 - Dealing with failures
 - Performance & scalability
 - Usability
-

What is MapReduce?



A programming model for large-scale distributed data processing

- Simple, elegant concept
- Restricted, yet powerful programming construct
- Building block for other parallel programming tools
- Extensible for different applications

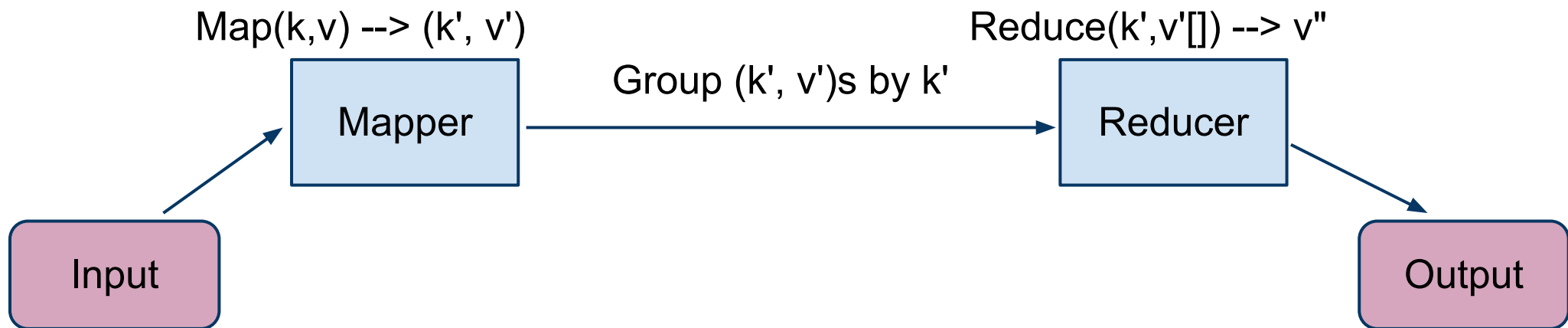
Also an implementation of a system to execute such programs

- Take advantage of parallelism
 - Tolerate failures and jitters
 - Hide messy internals from users
 - Provide tuning knobs for different applications
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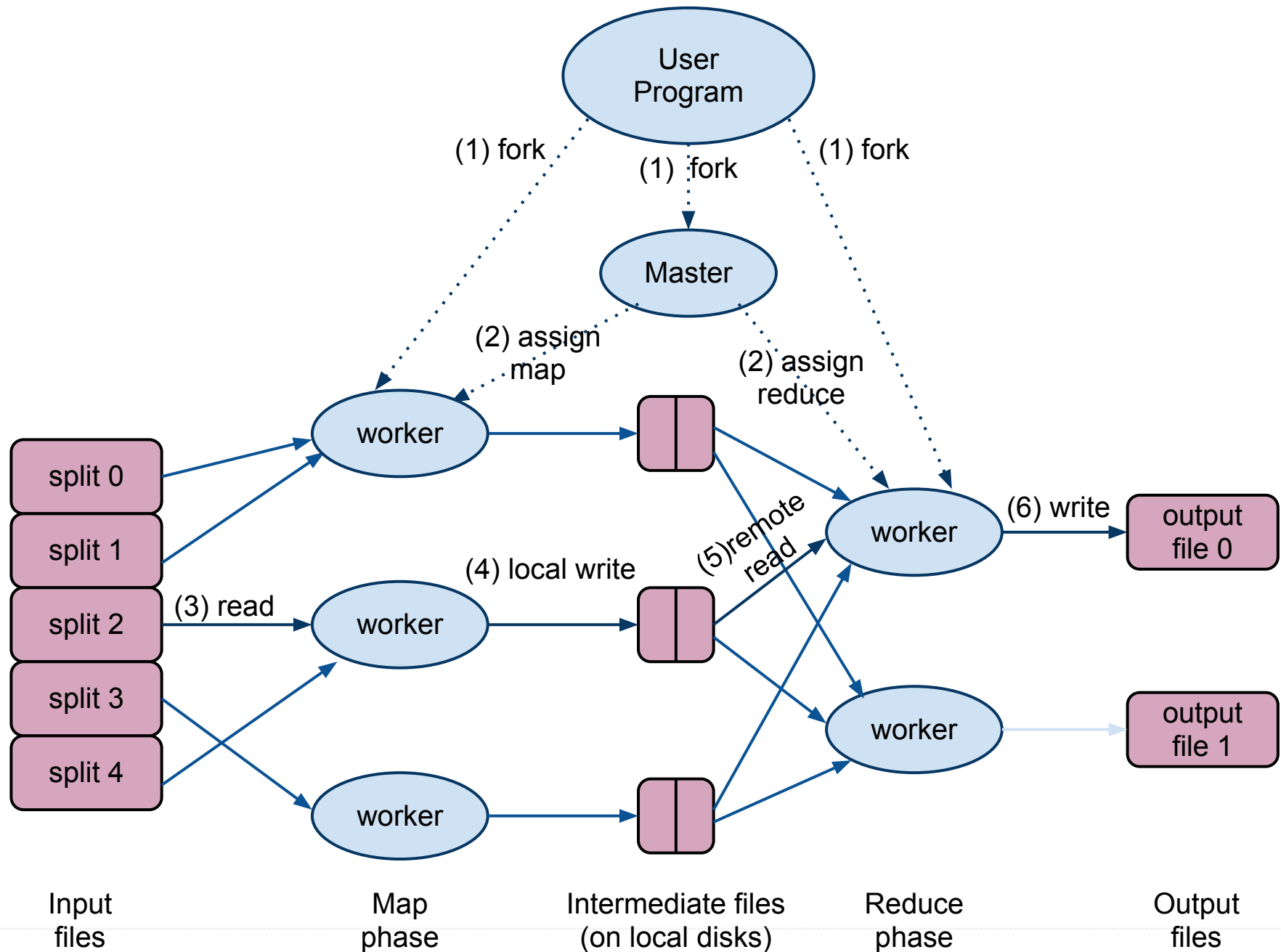
Programming Model



Inspired by Map/Reduce in functional programming languages, such as LISP from 1960's, but not equivalent



MapReduce Execution Overview



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Use of MapReduce inside Google



Stats for Month	Aug.'04	Mar.'06	Sep.'07
Number of jobs	29,000	171,000	2,217,000
Avg. completion time (secs)	634	874	395
Machine years used	217	2,002	11,081
Map input data (TB)	3,288	52,254	403,152
Map output data (TB)	758	6,743	34,774
reduce output data (TB)	193	2,970	14,018
Avg. machines per job	157	268	394
Unique implementations			
Mapper	395	1958	4083
Reducer	269	1208	2418

[From "MapReduce: simplified data processing on large clusters"](#)

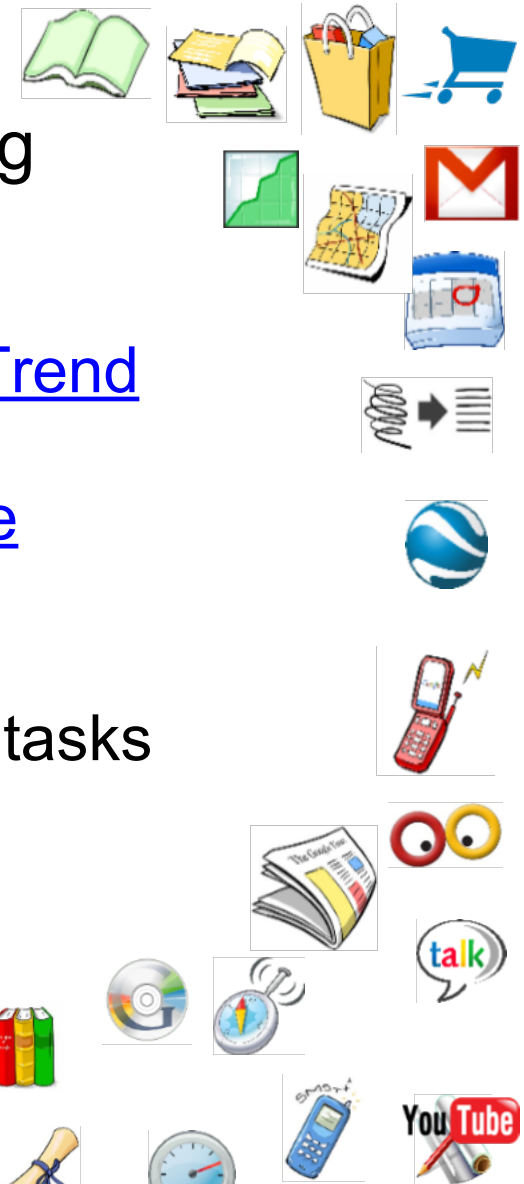
MapReduce inside Google



Googlers' hammer for 80% of our data crunching

- [Large-scale web search indexing](#)
- Clustering problems for [Google News](#)
- Produce reports for popular queries, e.g. [Google Trend](#)
- Processing of [satellite imagery data](#)
- Language model processing for [statistical machine translation](#)
- Large-scale [machine learning problems](#)
- Just a plain tool to reliably spawn large number of tasks
 - e.g. parallel data backup and restore

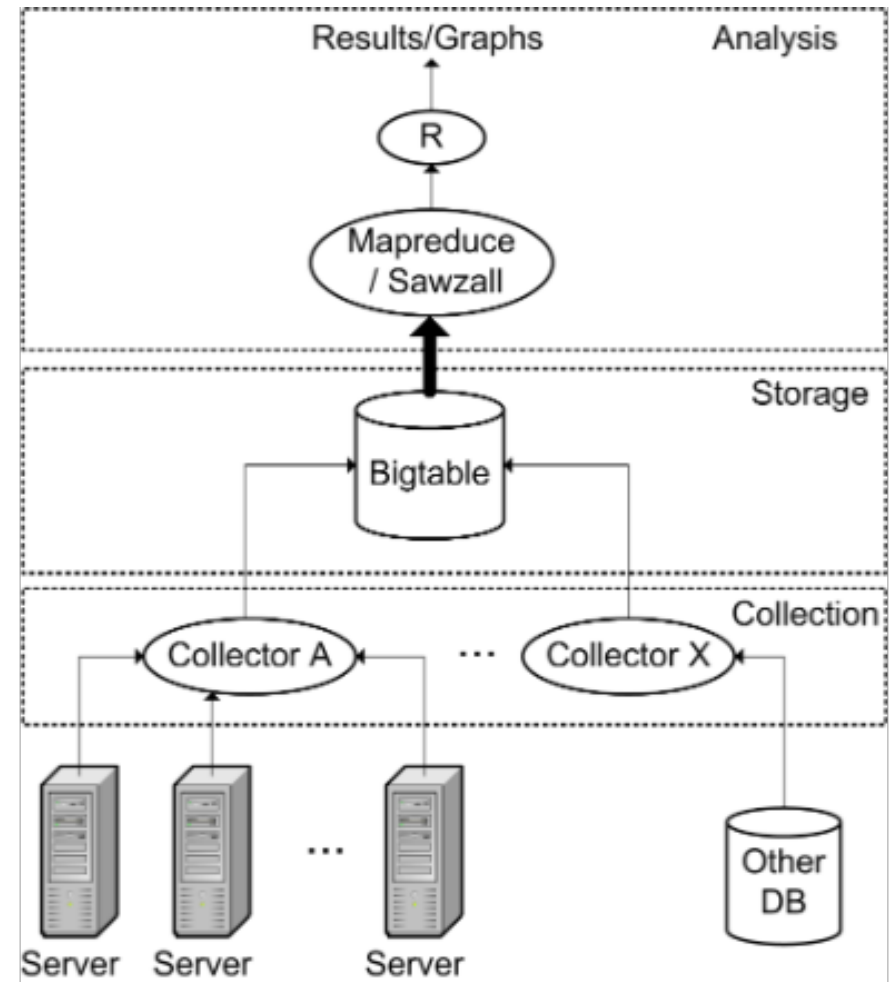
The other 20%? e.g. [Pregel](#)



Use of MR in System Health Monitoring



- Monitoring service talks to every server frequently
- Collect
 - Health signals
 - Activity information
 - Configuration data
- Store time-series data forever
- Parallel analysis of repository data
 - MapReduce/Sawzall



- Case study
 - Higher DRAM errors observed in a new GMail cluster
 - Similar servers running GMail elsewhere not affected
 - Same version of the software, kernel, firmware, etc.
 - Bad DRAM is the initial culprit
 - ... but that same DRAM model was fairly healthy elsewhere
 - Actual problem: bad motherboard batch
 - Poor electrical margin in some memory bus signals
 - GMail got more than its fair share of the bad batch
 - Analysis of this batch allocated to other services confirmed the theory

 - Analysis possible by having all relevant data in one place and processing power to digest it
 - MapReduce is part of the infrastructure
-

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- Word count and frequency in a large set of documents
 - Power of sorted keys and values
 - Combiners for map output
- Computing average income in a city for a given year
 - Using customized readers to
 - Optimize MapReduce
 - Mimic **rudimentary** DBMS functionality
- Overlaying satellite images
 - Handling various input formats using protocol buffers

Word Count Example



- Input: Large number of text documents
- Task: Compute word count across all the documents

Solution

- Mapper:
 - For every word in a document output (word, "1")
- Reducer:
 - Sum all occurrences of words and output (word, total_count)

Word Count Solution



```
//Pseudo-code for "word counting"
map(String key, String value):
  // key: document name,
  // value: document contents
  for each word w in value:
    EmitIntermediate(w, "1");

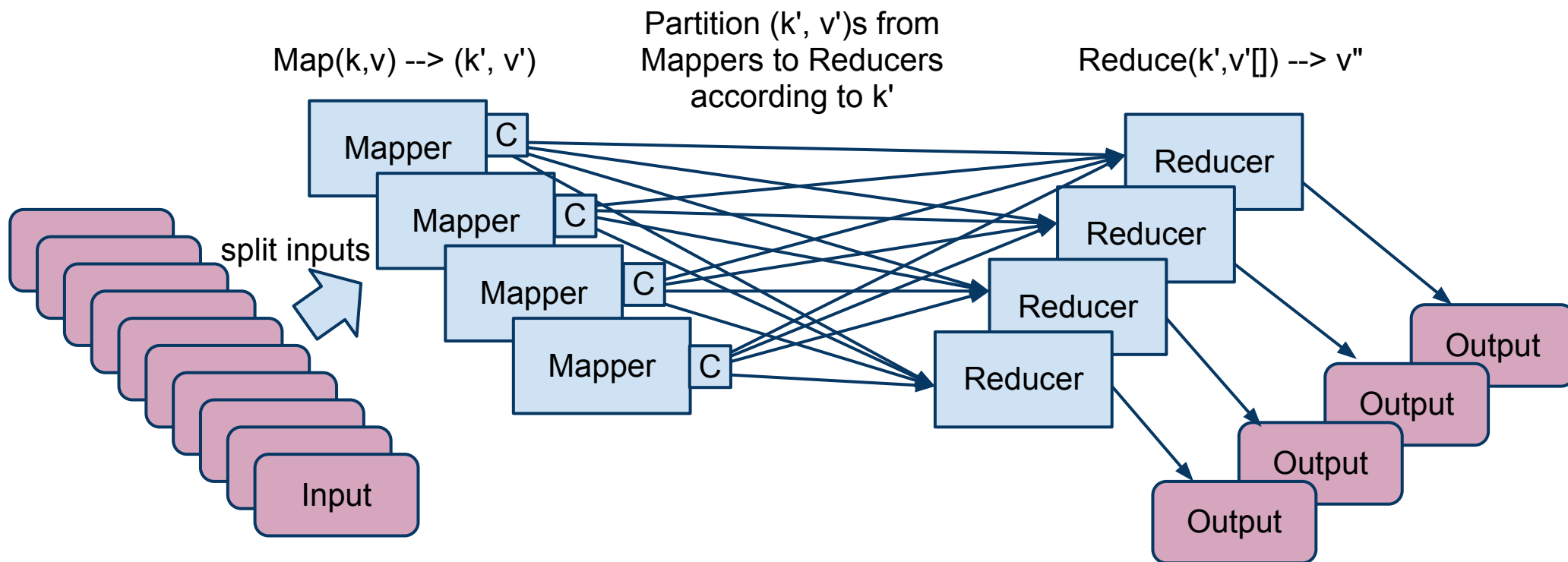
reduce(String key, Iterator values):
  // key: a word
  // values: a list of counts
  int word_count = 0;
  for each v in values:
    word_count += ParseInt(v);
  Emit(key, AsString(word_count));
```

No types, just strings*

Word Count Optimization: Combiner



- Apply reduce function to map output before it is sent to reducer
 - Reduces number of records output by the mapper!



Word Probability Example



- Input: Large number of text documents
- Task: Compute word probabilities across all the documents
 - Frequency is calculated using the total word count
- A naive solution with basic MapReduce model requires two MapReduces
 - MR1: count number of all words in these documents
 - Use combiners
 - MR2: count number of each word and divide it by the total count from MR1

- Can we do better?
- Two nice features of Google's MapReduce implementation
 - Ordering guarantee of reduce key
 - Auxiliary functionality: `EmitToAllReducers(k, v)`
- A nice trick: To compute the total number of words in all documents
 - Every map task sends its total word count with key "" to ALL reducer splits
 - Key "" will be the first key processed by reducer
 - Sum of its values → total number of words!

Word Probability Solution: Mapper with Combiner



```
map(String key, String value) :  
    // key: document name, value: document contents  
    int word_count = 0;  
    for each word w in value:  
        EmitIntermediate(w, "1");  
        word_count++;  
    EmitIntermediateToAllReducers("", AsString(word_count));  
  
combine(String key, Iterator values) :  
    // Combiner for map output  
    // key: a word, values: a list of counts  
    int partial_word_count = 0;  
    for each v in values:  
        partial_word_count += ParseInt(v);  
    Emit(key, AsString(partial_word_count));
```

Word Probability Solution: Reducer



```
reduce(String key, Iterator values):
    // Actual reducer
    // key: a word
    // values: a list of counts
    if (is_first_key):
        assert("" == key); // sanity check
        total_word_count_ = 0;
        for each v in values:
            total_word_count_ += ParseInt(v)
    else:
        assert("" != key); // sanity check
        int word_count = 0;
        for each v in values:
            word_count += ParseInt(v);
        Emit(key, AsString(word_count / total_word_count_));
```

- Word frequency in a large set of documents
 - Power of sorted keys and values
 - Combiners for map output
- **Computing average income in a city for a given year**
 - Using customized readers to
 - Optimize MapReduce
 - Mimic **rudimentary** DBMS functionality
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Average Income In a City



SSTable 1: (SSN, {Personal Information})

123456:(John Smith;Sunnyvale, CA)

123457:(Jane Brown;Mountain View, CA)

123458:(Tom Little;Mountain View, CA)

SSTable 2: (SSN, {year, income})

123456:(2007,\$70000),(2006,\$65000),(2005,\$6000),...

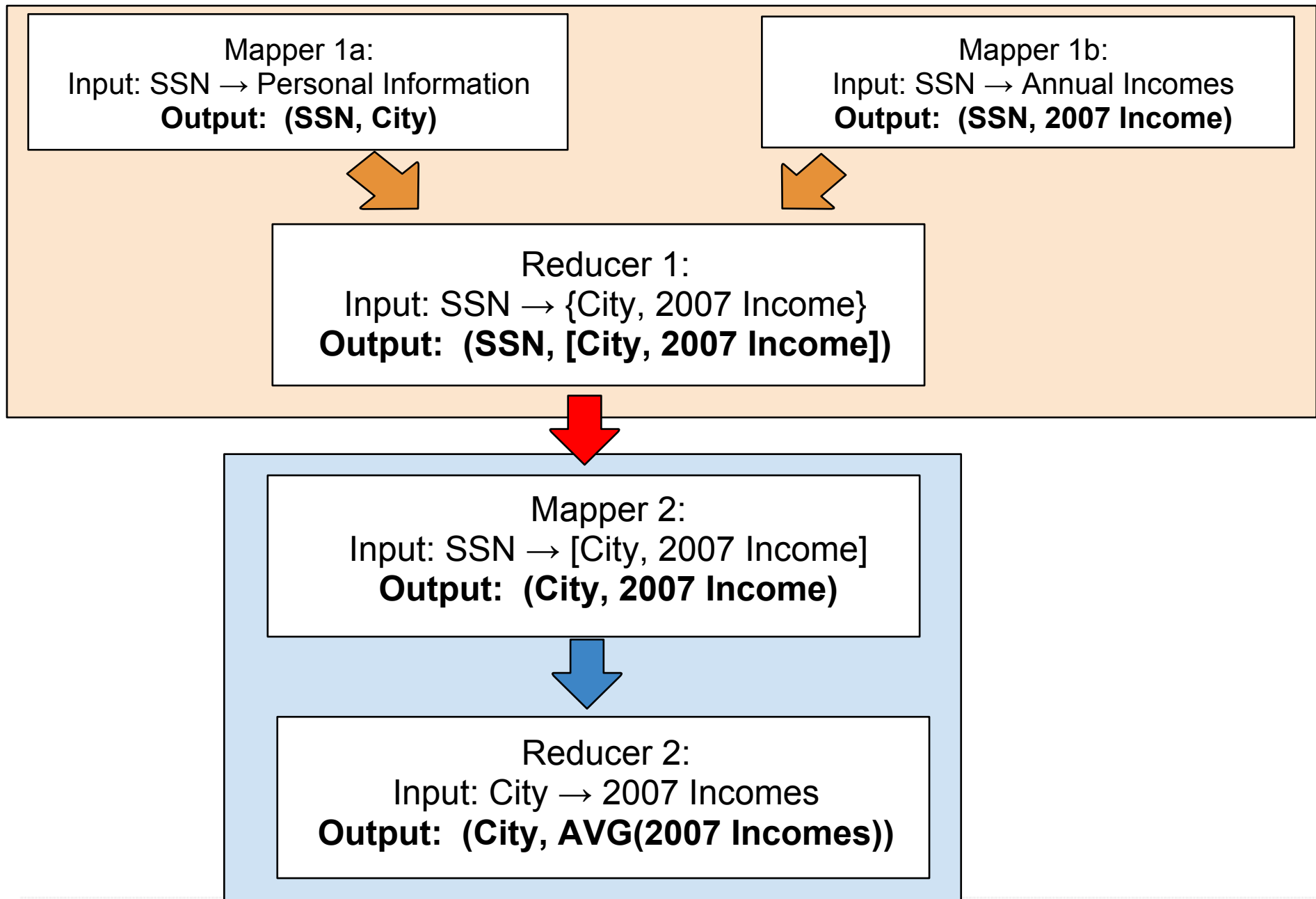
123457:(2007,\$72000),(2006,\$70000),(2005,\$6000),...

123458:(2007,\$80000),(2006,\$85000),(2005,\$7500),...

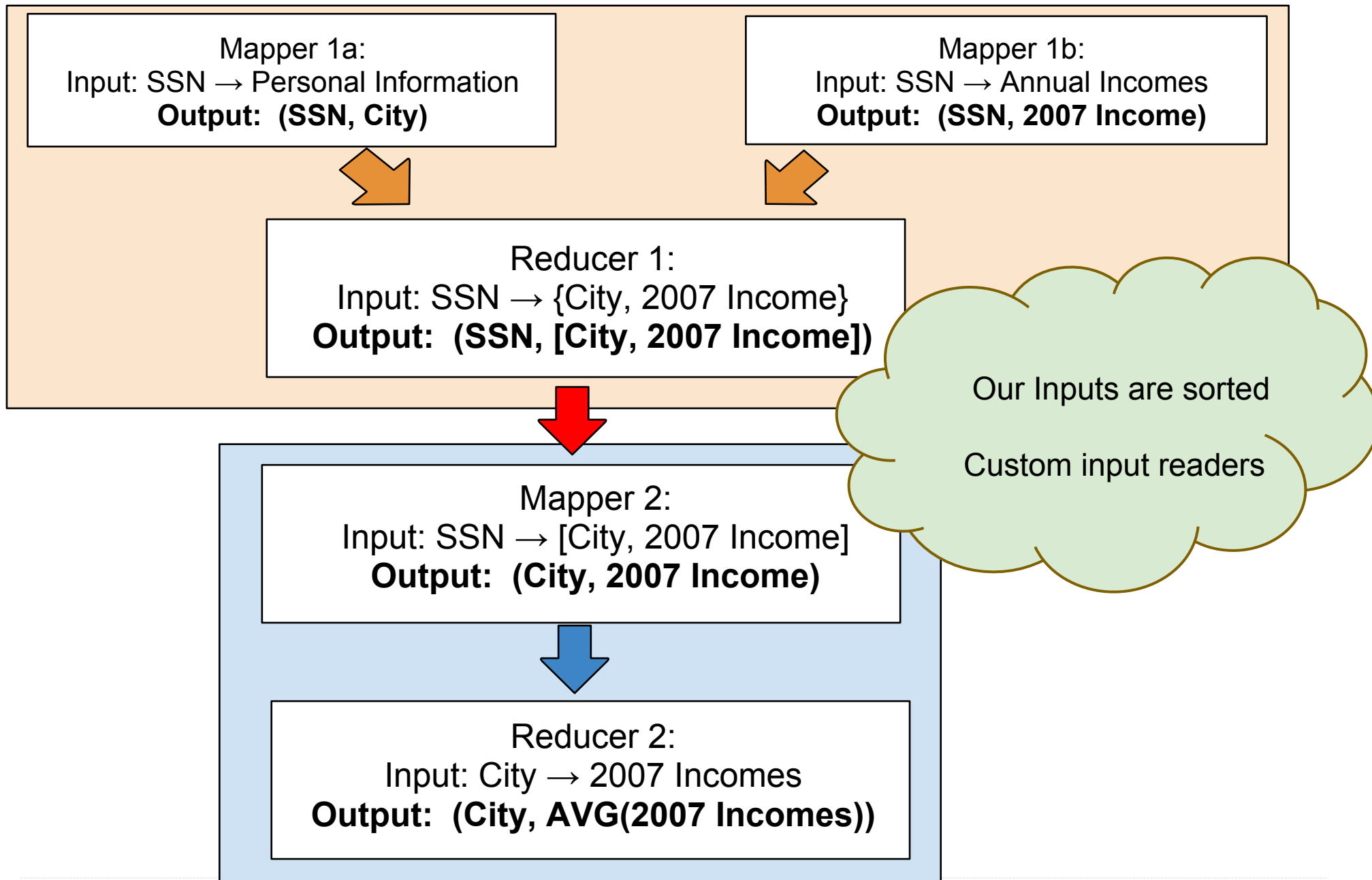
Task: Compute average income in each city in 2007

Note: **Both inputs sorted by SSN**

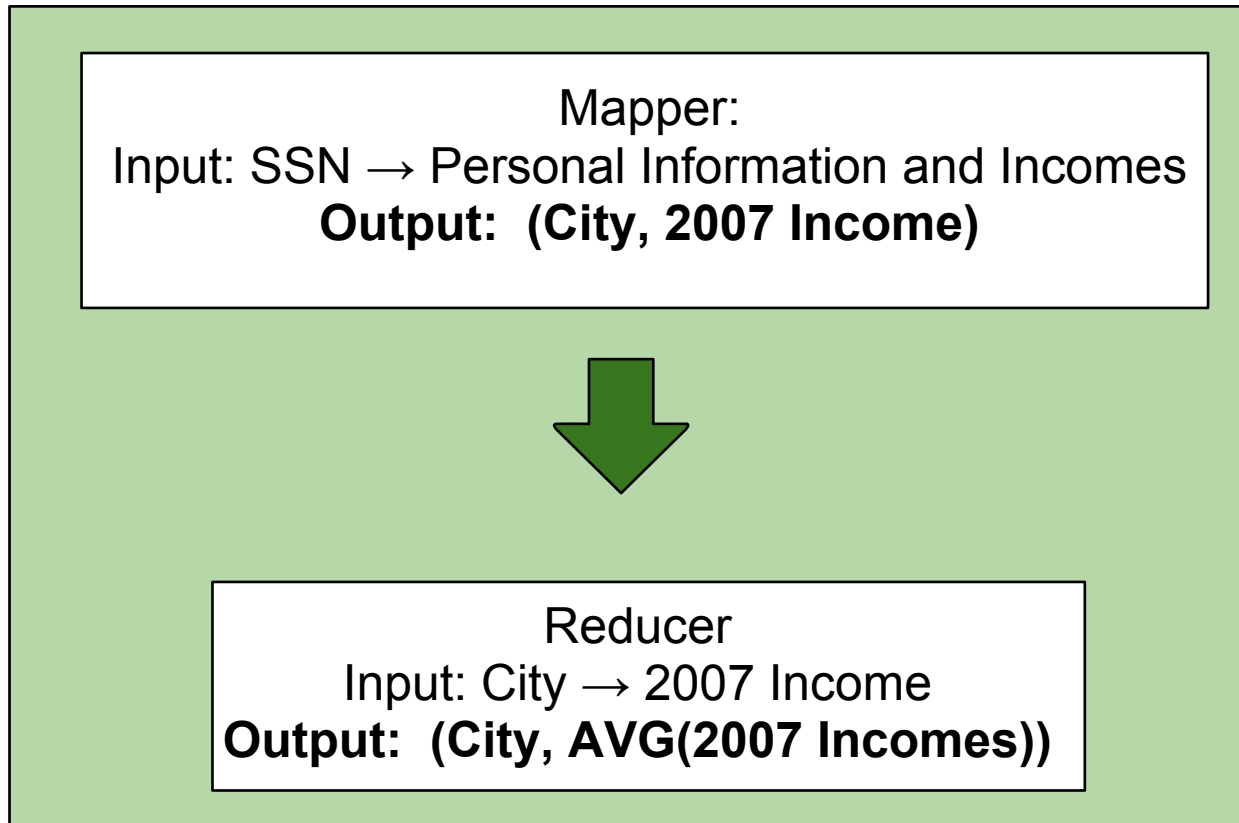
Average Income in a City Basic Solution



Average Income in a City Basic Solution



Average Income in a Joined Solution



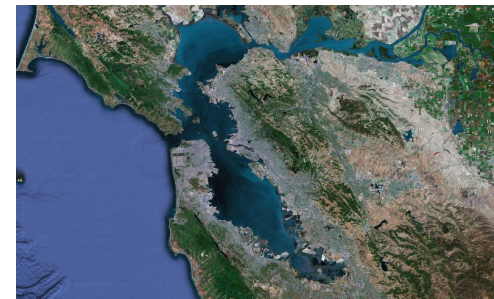
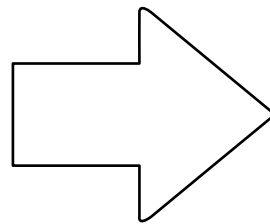
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Stitch Imagery Data for Google Maps



A simplified version could be:

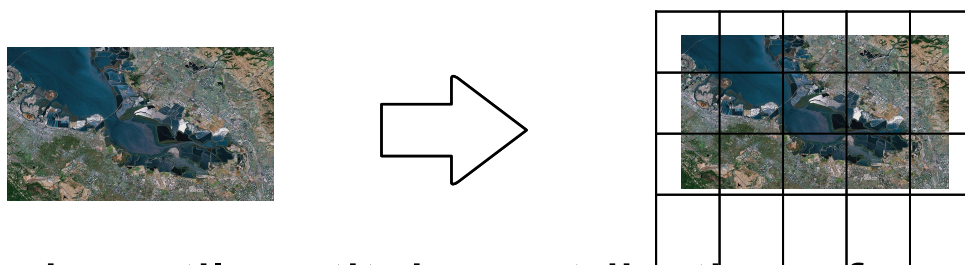
- Imagery data from different content providers
 - Different formats
 - Different coverages
 - Different timestamps
 - Different resolutions
 - Different exposures/tones
- Large amount to data to be processed
- Goal: produce data to serve a "satellite" view to users



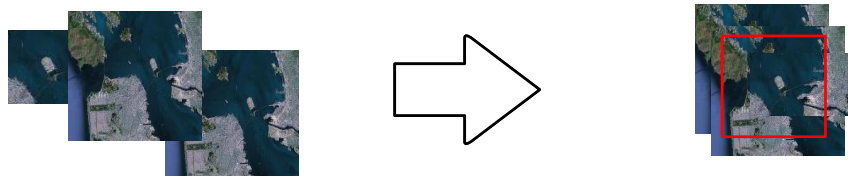
Stitch Imagery Data Algorithm



1. Split the whole territory into "tiles" with fixed location IDs
2. Split each source image according to the tiles it covers



3. For a given tile, stitch contributions from different sources, based on its freshness and resolution, or other preference



4. Serve the merged imagery data for each tile, so they can be loaded into and served from an image server farm.

Using Protocol Buffers to Encode Structured Data



- Open sourced from Google, among [many others](http://code.google.com/p/protobuf/):
<http://code.google.com/p/protobuf/>
- It supports C++, Java and Python.
- A way of encoding structured data in an efficient yet extensible format. e.g. we can define

```
message Tile {
  required int64 location_id = 1;
  group coverage {
    double latitude = 2;
    double longitude = 3;
    double width = 4; // in km
    double length = 5; // in km
  }
  required bytes image_data = 6; // Bitmap Image data
  required int64 timestamp = 7;
  optional float resolution = 8 [default = 10];
  optional string debug_info = 10;
}
```

Google uses Protocol Buffers for almost all its internal RPC protocols, file formats and of course in MapReduce.

Stitch Imagery Data Solution: Mapper



```
map(String key, String value) :
  // key: image file name
  // value: image data
  Tile whole_image;

  switch (file_type(key)) :
    FROM_PROVIDER_A: Convert_A(value, &whole_image);
    FROM_PROVIDER_B: Convert_B(...);
    ...

  // split whole_image according to the grid into tiles
  for each Tile t in whole_image
    string v;
    t.SerializeToString(&v);
    EmitIntermediate(IntToStr(t.location_id()), v);
```

Stitch Imagery Data Solution: Reducer



```
reduce(String key, Iterator values):  
    // key: location_id,  
    // values: tiles from different sources  
  
    sort values according v.resolution() and v.timestamp();  
  
    Tile merged_tile;  
    for each v in values:  
        overlay pixels in v to merged_tile based on  
            v.coverage();  
  
    Normalize merged_tile to be the serve tile size;  
  
    Emit(key, ProtobufToString(merged_tile));
```

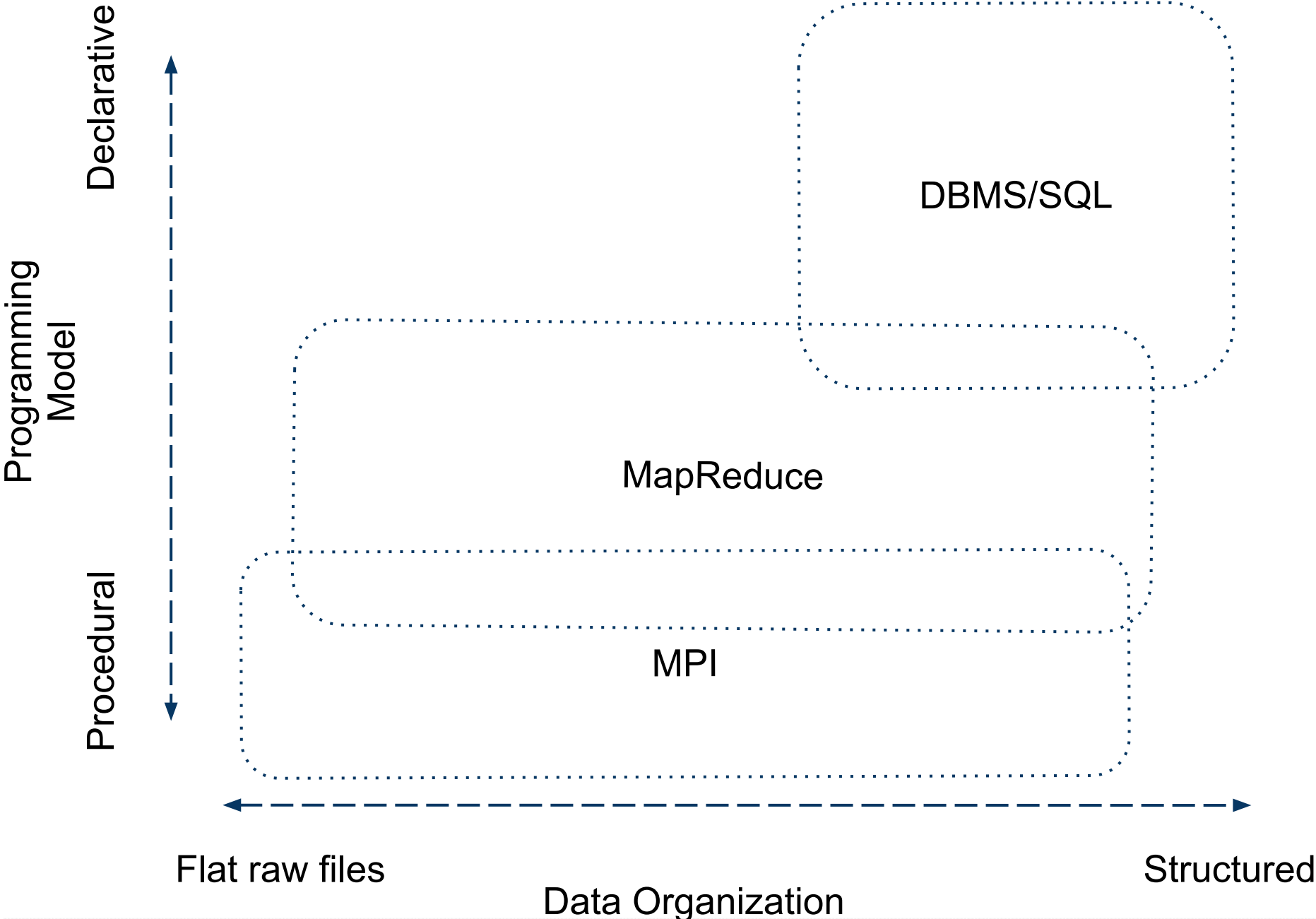
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Dimensions to compare Apples and Oranges

- Data organization
- Programming model
- Execution model

- Target applications
- Assumed computing environment
- Overall operating cost

My Basket of Fruit



Nutritional Information of My Basket



	MPI	MapReduce	DBMS/SQL
What they are	A general parallel programming paradigm	A programming paradigm and its associated execution system	A system to store, manipulate and serve data.
Programming Model	Messages passing between nodes	Restricted to Map/Reduce operations	Declarative on data query/retrieving; Stored procedures
Data organization	No assumption	"files" can be sharded	Organized datastructures
Data to be manipulated	Any	k,v pairs: string/ protomsg	Tables with rich types
Execution model	Nodes are independent	Map/Shuffle/Reduce Checkpointing/Backup Physical data locality	Transaction Query/operation optimization Materialized view
Usability	Steep learning curve*; difficult to debug	Simple concept Could be hard to optimize	Declarative interface; Could be hard to debug in runtime
Key selling point	Flexible to accommodate various applications	Plow through large amount of data with commodity hardware	Interactive querying the data; Maintain a consistent view across clients

See what others say: [\[1\]](#), [\[2\]](#), [\[3\]](#)

Dimensions to choose between Apples and Oranges for an application developer:

- Target applications
 - Complex operations run frequently v.s. one time plow
 - Off-line processing v.s. real-time serving
 - Assumed computing environment
 - Off-the-shelf, custom-made or [donated](#)
 - Formats and sources of your data
 - Overall operating cost
 - Hardware maintenance, license fee
 - Manpower to develop, monitor and debug
-

Existing MapReduce and Similar Systems



Google MapReduce

- Support C++, Java, Python, Sawzall, etc.
- Based on proprietary infrastructures
 - [GFS\(SOSP'03\)](#), [MapReduce\(OSDI'04\)](#) , [Sawzall\(SPJ'05\)](#), [Chubby\(OSDI'06\)](#), [Bigtable\(OSDI'06\)](#)
 - and some [open source libraries](#)

Hadoop Map-Reduce

- Open Source!
- Plus the whole equivalent package, and more
 - HDFS, Map-Reduce, Pig, Zookeeper, HBase, Hive
- Used by Yahoo!, Facebook, Amazon and [Google-IBM NSF cluster](#)

Dryad

- Proprietary, based on Microsoft SQL servers
- [Dryad\(EuroSys'07\)](#), [DryadLINQ\(OSDI'08\)](#)
- [Michael's Dryad TechTalk@Google \(Nov.'07\)](#)

And [others](#)

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Inverted Index Construction



- Input: Large number of text documents
- Task: Postings lists for every term in the collection
 - For every word, all documents that contain the word and the positions.



Solution:

- Mapper:
 - For every word in a document output (word, [URL, position])
- Reducer:
 - Aggregate all the information that we have about each word.

Inverted Index Solution



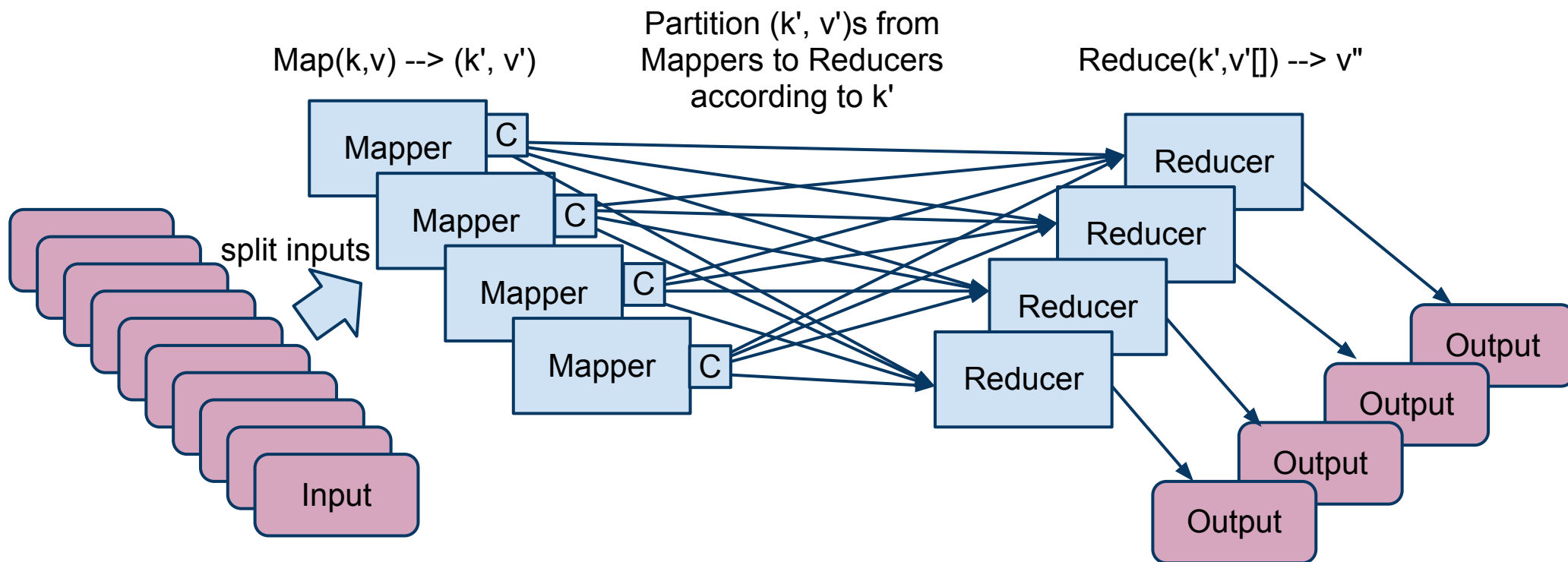
```
//Pseudo-code for "inverted index"
map(String key, String value):
  // key: document URL,
  // value: document contents
  vector words = tokenize(value)
  for position from 0 to len(words):
    EmitIntermediate(w, {key, position});

reduce(String key, Iterator values):
  // key: a word
  // values: a list of {URL, position} tuples.
  postings_list = [];
  for each v in values:
    postings_list.append(v);
  sort(postings_list); // Sort by URL, then position
  Emit(key, AsString(postings_list));
```

Inverted Index Optimization: Combiner



- Combiners can also be used to reduce the number of intermediate outputs, to start aggregating all occurrences of document words.



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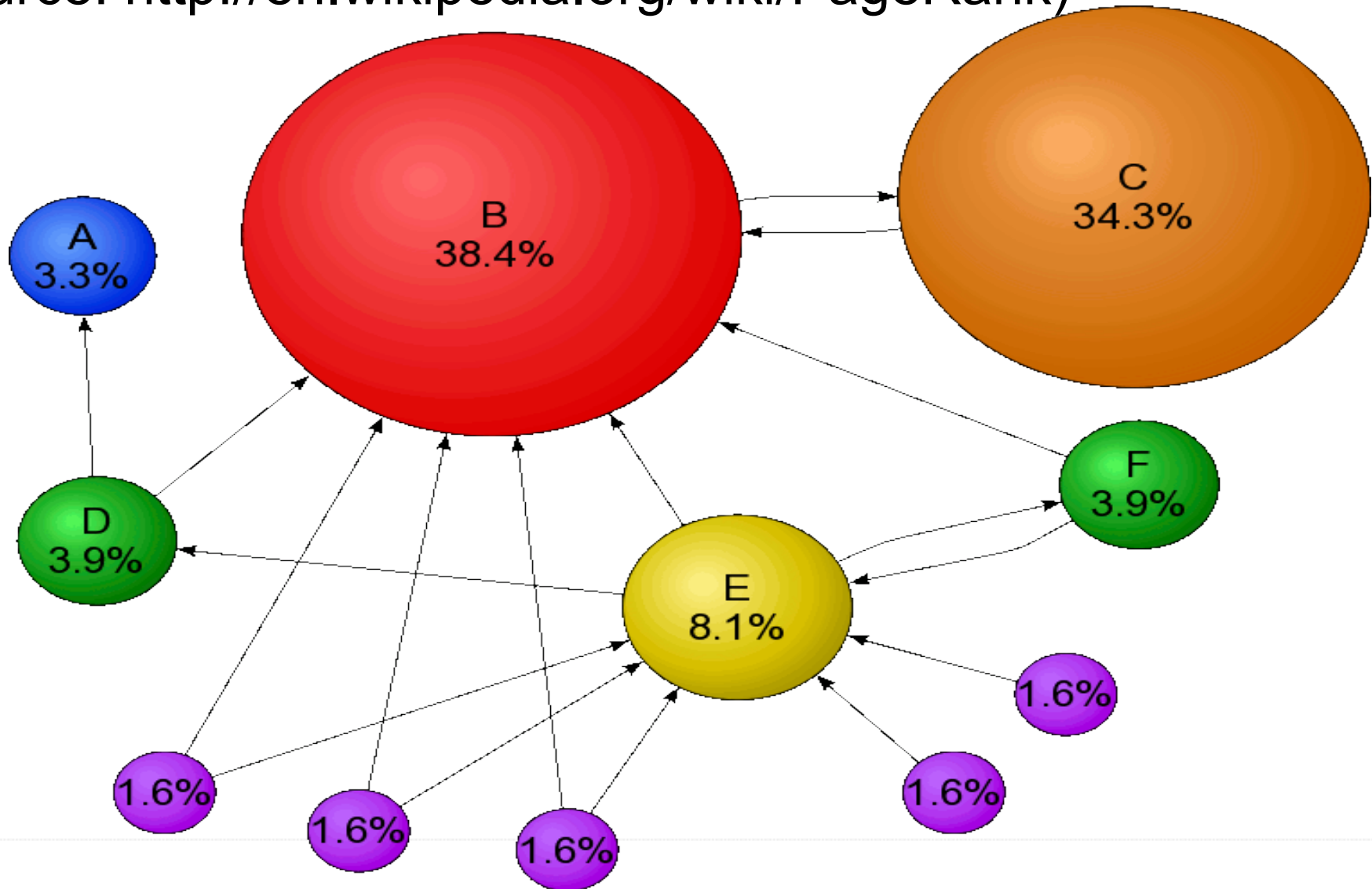
 - Implementation of Google MapReduce
 - Dealing with failures
 - Performance & scalability
 - (Operational) Usability
 - monitoring, debugging, profiling, etc.

 - Practical indexing examples in IR
 - Inverted index construction
 - **PageRank computation**
-

- Input: Large number of documents with hyperlinks structured as a graph.
 - Task:
 - Algorithm to compute the probability that a random walk on the graph will land in a given page.
 - Used as a measure of the importance of the page.
 - With a small probability, the user can jump to any page in the graph (not following hyperlinks).
-

PageRank computation

(Source: <http://en.wikipedia.org/wiki/PageRank>)

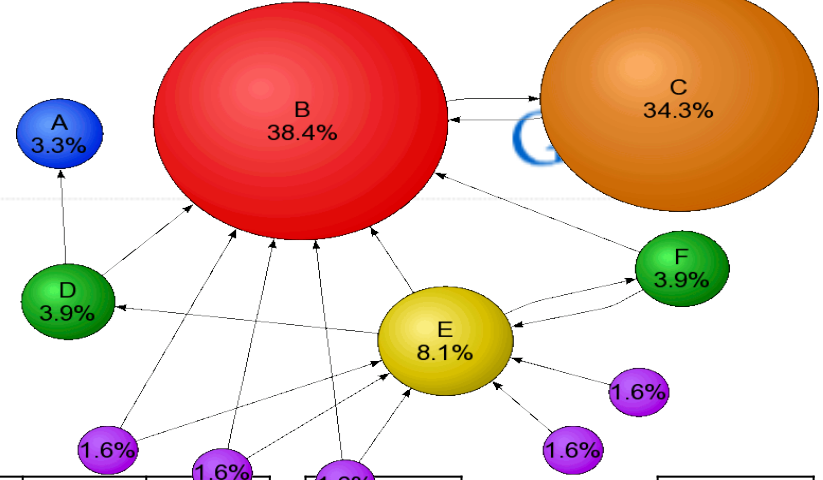


Algorithm:

- N = total number of web pages
- Matrix M defined as;
 - $M[i][j]$ is 0 if the j -th page has no links to the i -th page.
 - $M[i][j]$ is the probability to move from page j to page i , assuming the same probability for all outgoing links.
- Vector R defined as:
 - $R[i]$ is the estimated PageRank value for page i .
- Iterative algorithm:
$$R = (1-d) \cdot \mathbf{1} + d \cdot M \cdot R$$
where d is the decay term

PageRank computation

No decay term, one iteration.
Most probability mass in B and E

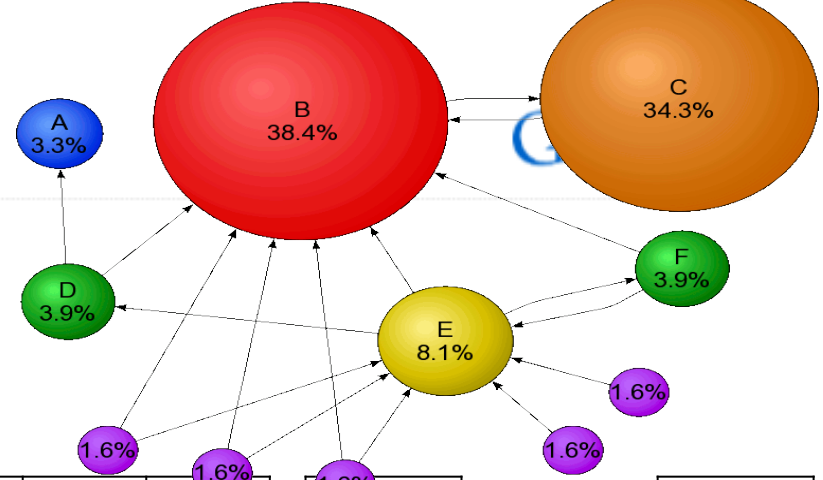


	A	B	C	D	E	F	G	H	I	J	K	PR	PR
A	0	0	0	0.5	0	0	0	0	0	0	0	0.09	0.05
B	0	0	1	0.5	0.33	0.5	0.5	0.5	0.5	0	0	0.09	0.34
C	0	1	0	0	0	0	0	0	0	0	0	0.09	0.09
D	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.03
E	0	0	0	0	0	0.5	0.5	0.5	0.5	1	1	0.09	0.36
F	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.03
G	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
H	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
I	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
J	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
K	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00

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PageRank computation

No decay term, three iterations.
Most probability mass in B and C

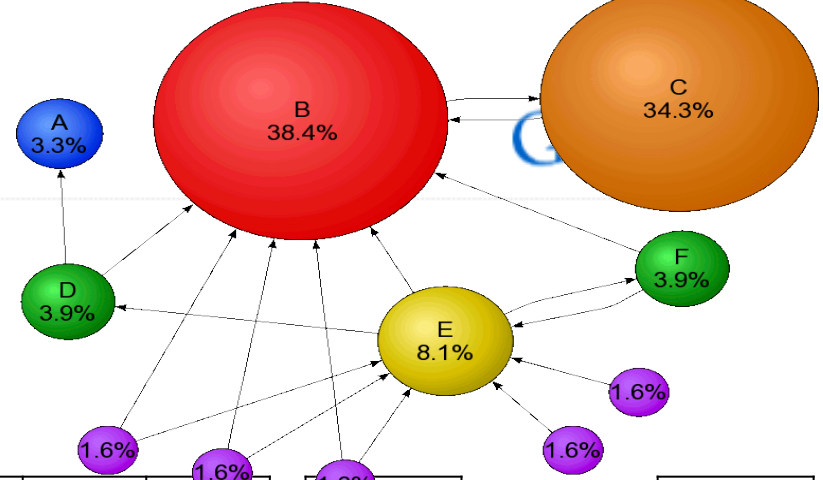


	A	B	C	D	E	F	G	H	I	J	K	PR	PR
A	0	0	0	0.5	0	0	0	0	0	0	0	0.09	0.06
B	0	0	1	0.5	0.33	0.5	0.5	0.5	0.5	0	0	0.09	0.47
C	0	1	0	0	0	0	0	0	0	0	0	0.09	0.24
D	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.01
E	0	0	0	0	0	0.5	0.5	0.5	0.5	1	1	0.09	0.06
F	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.01
G	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
H	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
I	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
J	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00
K	0	0	0	0	0	0	0	0	0	0	0	0.09	0.00

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PageRank computation

Decay term 0.18, three iterations.



	A	B	C	D	E	F	G	H	I	J	K	PR	PR
A	0	0	0	0.5	0	0	0	0	0	0	0	0.09	0.04
B	0	0	1	0.5	0.33	0.5	0.5	0.5	0.5	0	0	0.09	0.45
C	0	1	0	0	0	0	0	0	0	0	0	0.09	0.57
D	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.06
E	0	0	0	0	0	0.5	0.5	0.5	0.5	1	1	0.09	0.09
F	0	0	0	0	0.33	0	0	0	0	0	0	0.09	0.01
G	0	0	0	0	0	0	0	0	0	0	0	0.09	0.01
H	0	0	0	0	0	0	0	0	0	0	0	0.09	0.01
I	0	0	0	0	0	0	0	0	0	0	0	0.09	0.01
J	0	0	0	0	0	0	0	0	0	0	0	0.09	0.01
K	0	0	0	0	0	0	0	0	0	0	0	0.09	0.01

=

- Matrix m is sparse:
 - We can store one $\langle \text{key}, \text{value} \rangle$ pair per row.
 - $\text{key} = \text{URL}$, $\text{value} = \text{URLs of outgoing links}$.
- Vector R :
 - one $\langle \text{key}, \text{value} \rangle$ pair per element.
- Matrix multiplication:
 - Join both sets (aggregate by key).
 - Multiply to produce each new value of R' in the reduce step.

- Joins are trivial to implement in MapReduce:
 - For the first dataset, one mapper function maps (key1, value1) to (key1, value1)
 - For the second dataset, other mapper function maps (key2, value2) to (key2, value2)
 - The reducer aggregates, for the same key, the two values, if both are present.

PageRank computation



```
//Pseudo-code for "PageRank" (no decay factor)
map_matrix(String key,
           String input_value,
           String joined_input_value):
    // key: document URL,
    // input_value: URLs of outgoing links and weights
    // joined_input_value: Current PageRank.
    for (URL, weight) in value:
        EmitIntermediate(key, (weight * joined_input_value));

reduce_pagerank(String key, Iterator values):
    // key: a URL
    // values: the incoming PageRank from each node with
    //          an incoming link.
    int sum = 0;
    for each v in values:
        sum += v;
    Emit(key, v);
```

PageRank computation



```
//Pseudo-code for "PageRank" (with decay factor)
map_matrix(String key,
           String input_value,
           String joined_input_value):
    // key: document URL,
    // input_value: URLs of outgoing links and weights
    // joined_input_value: Current PageRank.
    for (URL, weight) in value:
        EmitIntermediate(key, (weight * joined_input_value));

reduce_pagerank(String key, Iterator values):
    // key: a URL
    // values: the incoming PageRank from each node with
    //          an incoming link.
    int sum = 0;
    for each v in values:
        sum += v;

    // N is the graph size, assumed to be known from when the input
    // sparse matrix was constructed.
    sum = d * sum + (1-d) / N;
    Emit(key, v);
```

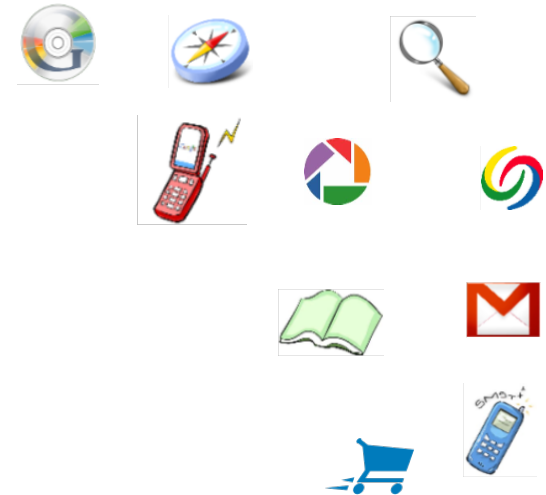
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 - Usability
-

Google Computing Infrastructure



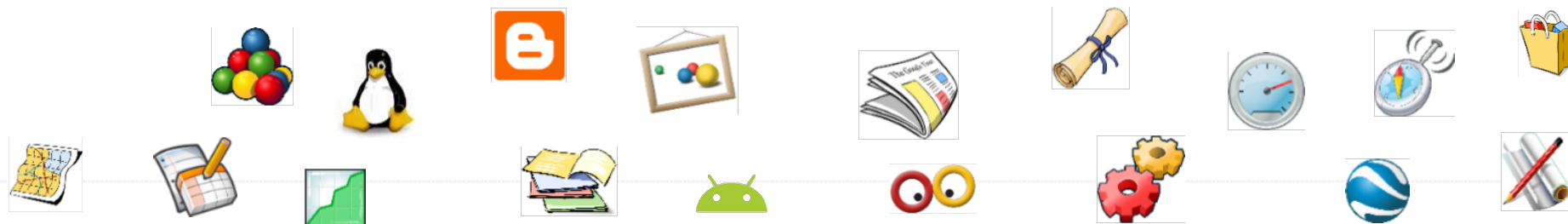
- Infrastructure must support

- Diverse set of applications
 - Increasing over time
- Ever-increasing application usage
- Ever-increasing computational requirements
- Cost effective

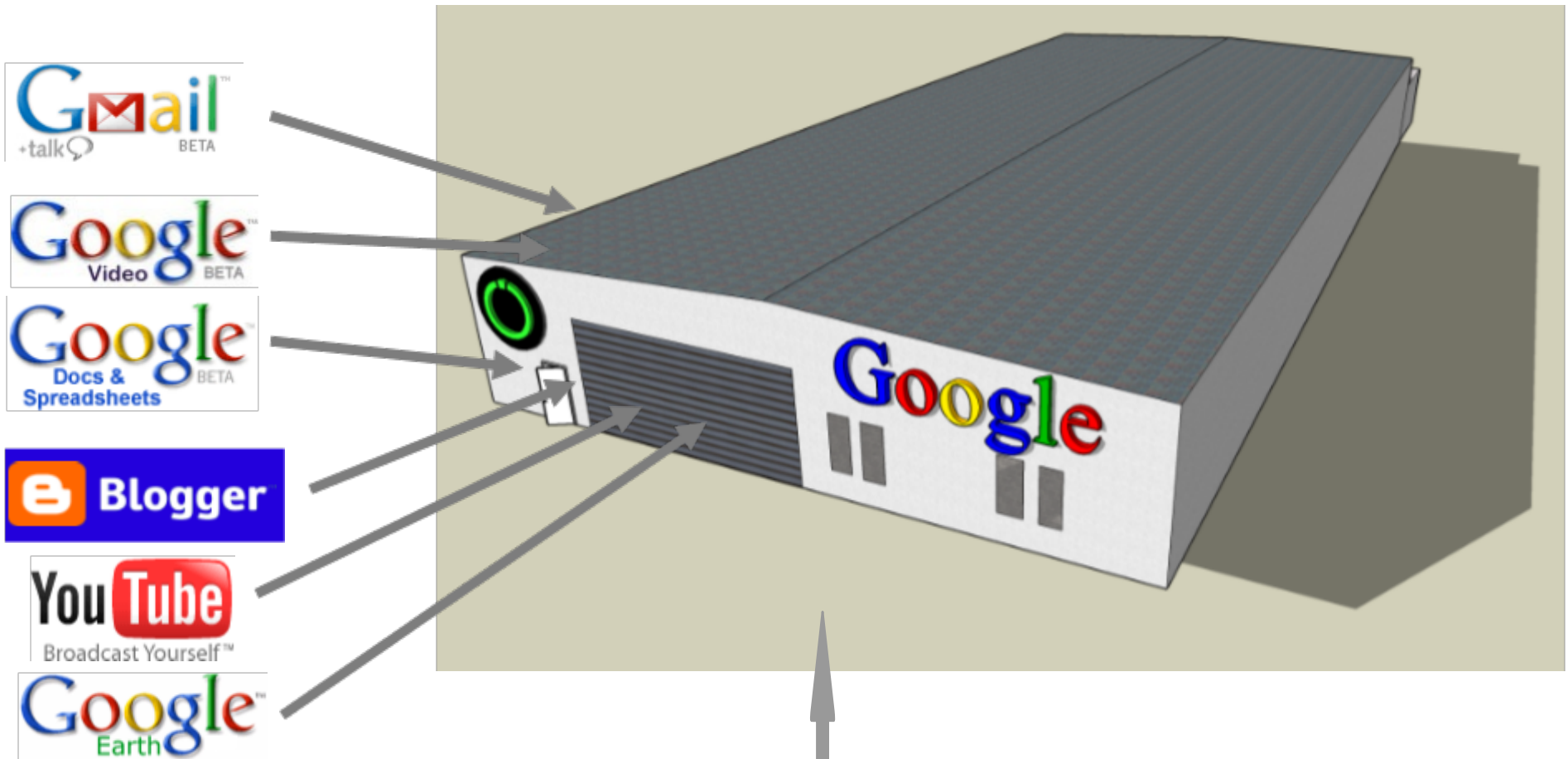


- Data centers

- Google-specific mechanical, [thermal](#) and [electrical](#) design
- Highly-customized PC-class motherboards
- Running Linux
- In-house management & application software



Sharing is the Way of Life



+ Batch processing
(MapReduce, Sazwall)

To organize the world's information and make it universally accessible and useful.

- Failure handling
 - Bad apples appear now and there
 - Scalability
 - Fast growing dataset
 - Broad extension of Google services
 - Performance and utilization
 - Minimizing run-time for individual jobs
 - Maximizing throughput across all services
 - Usability
 - Troubleshooting
 - Performance tuning
 - Production monitoring
-

- [LANL data](#) (DSN 2006)
 - Data collected over 9 years
 - Covered 4750 machines and 24101 CPUs
 - Distribution of failures
 - Hardware ~ 60%, Software ~ 20%, Network/Environment/Humans ~ 5%, Aliens ~ 25%*
 - Depending on a system, failures occurred between once a day to once a month
 - Most of the systems in the survey were the cream of the crop at their time
- [PlanetLab](#) (SIGMETRICS 2008 HotMetrics Workshop)
 - Average frequency of failures per node in a 3-months period
 - Hard failures: 2.1
 - Soft failures: 41
 - Approximately failure every 4 days

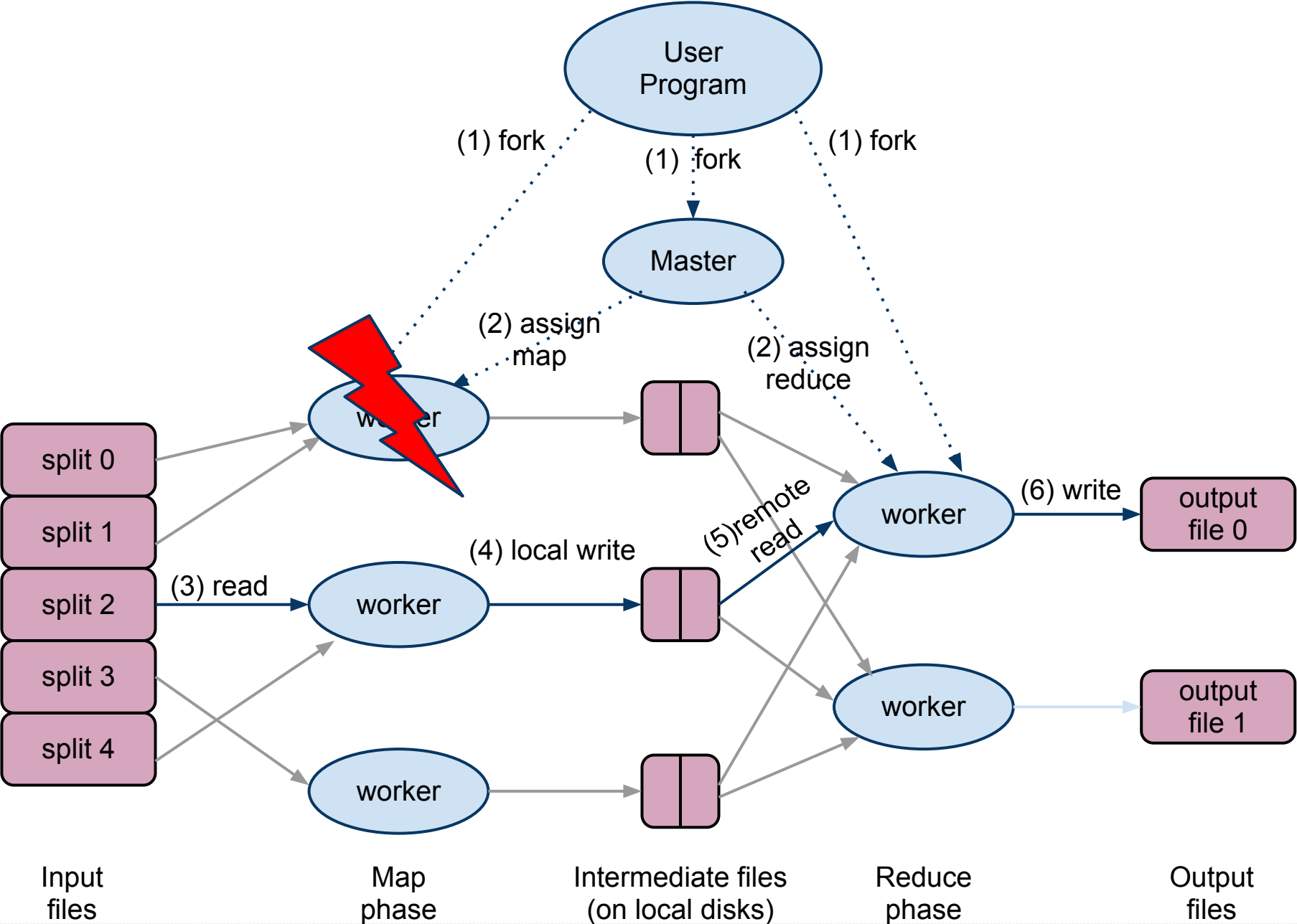
- [DRAM errors analysis](#) (SIGMETRICS 2009)
 - Data collected over 2.5 years
 - 25,000 to 70,000 errors per billion device hours per Mbit
 - Order of magnitude more than under lab conditions
 - 8% of DIMMs affected by errors
 - Hard errors are dominant cause of failure
 - [Disk drive failure analysis](#) (FAST 2007)
 - Annualized Failure Rates vary from 1.7% for one year old drives to over 8.6% in three year old ones
 - Utilization affects failure rates only in very old and very old disk drive populations
 - Temperature change can cause increase in failure rates but mostly for old drives
-

- Failures are a part of everyday life
 - Mostly due to the scale and shared environment
 - Sources of job failures
 - Hardware
 - Software
 - Preemption by a more important job
 - Unavailability of a resource due to overload
 - Failure types
 - Permanent
 - Transient
-

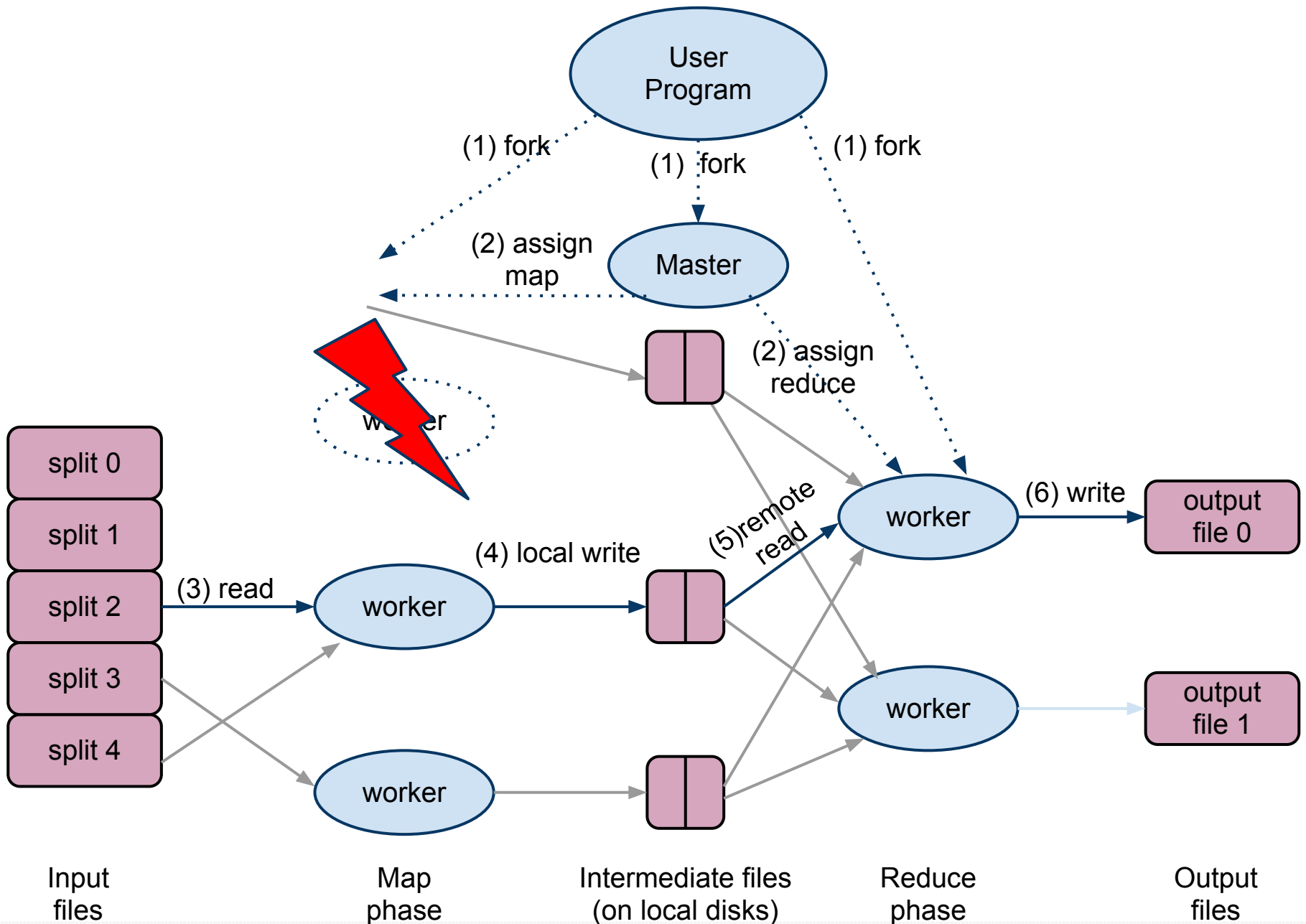
Different Failures Require Different Actions

- Fatal failure (the whole job dies)
 - Simplest case around :)
 - You'd prefer to resume computation rather than recompute
 - Transient failures
 - You'd want your job to adjust and finish when issues resolve
 - Program hangs.. forever.
 - Define "forever"
 - Can we figure out why?
 - What to do?
 - "It's-Not-My-Fault" failures
-

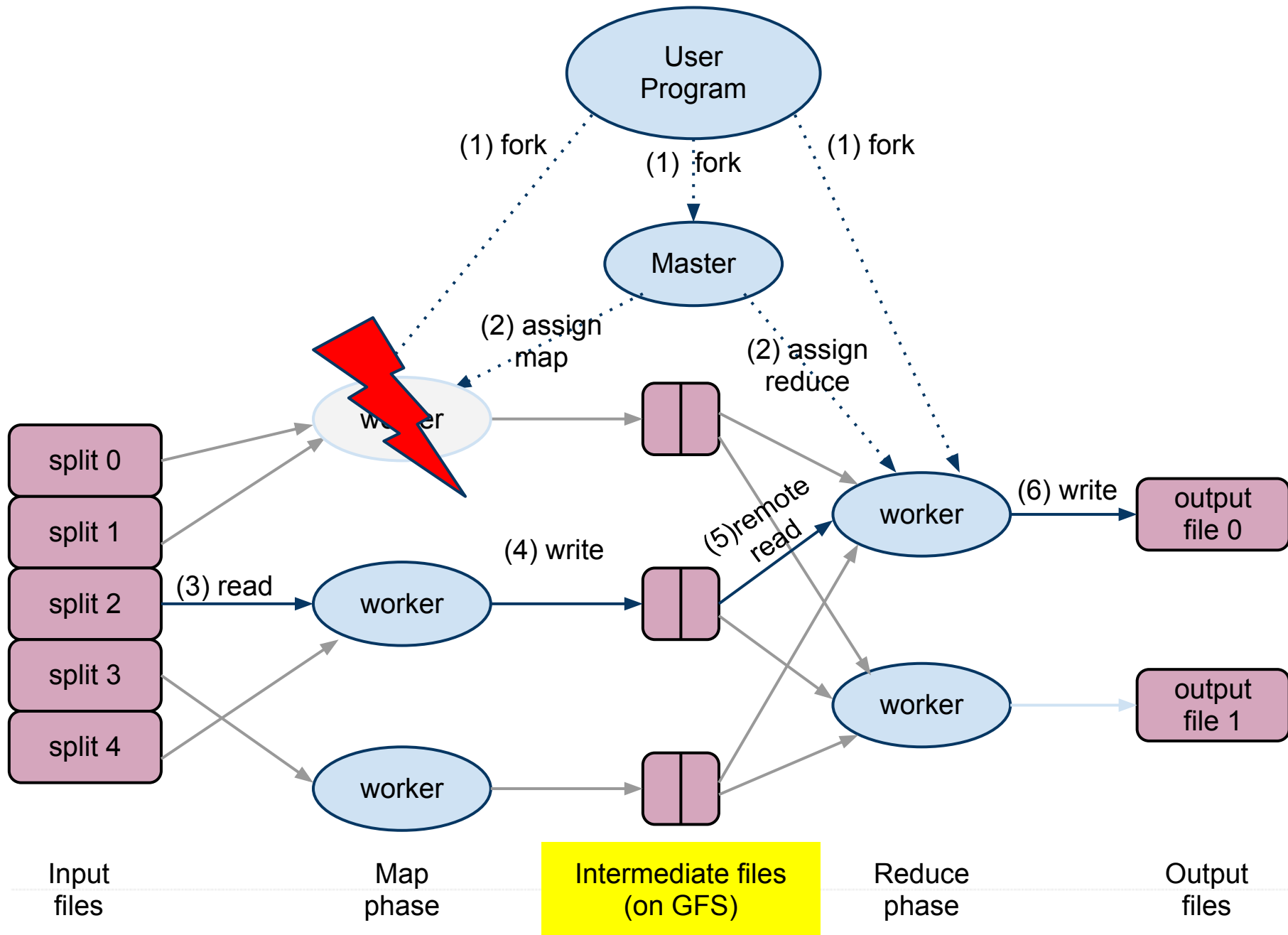
MapReduce: Task Failure



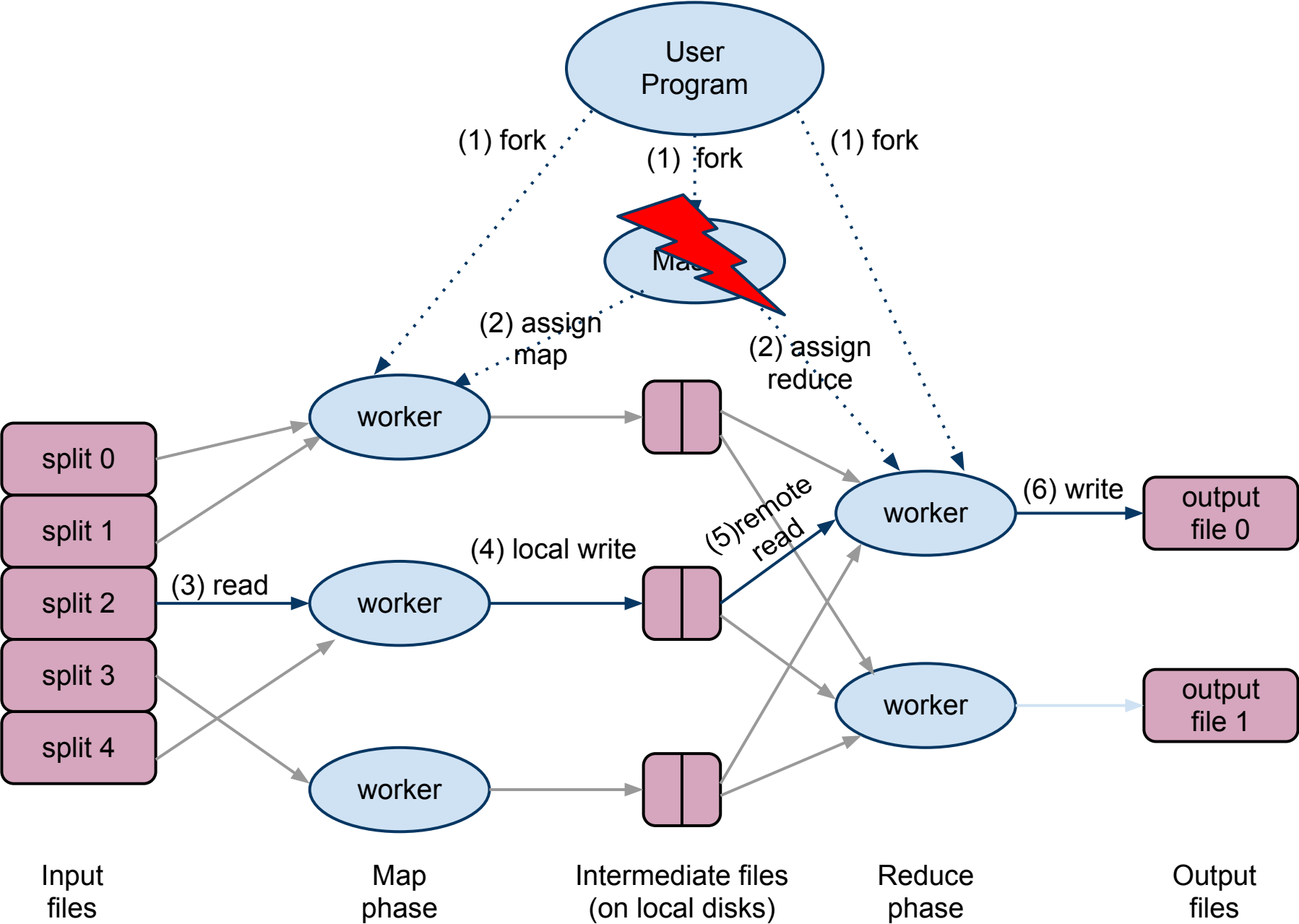
Recover from Task Failure by Re-execution



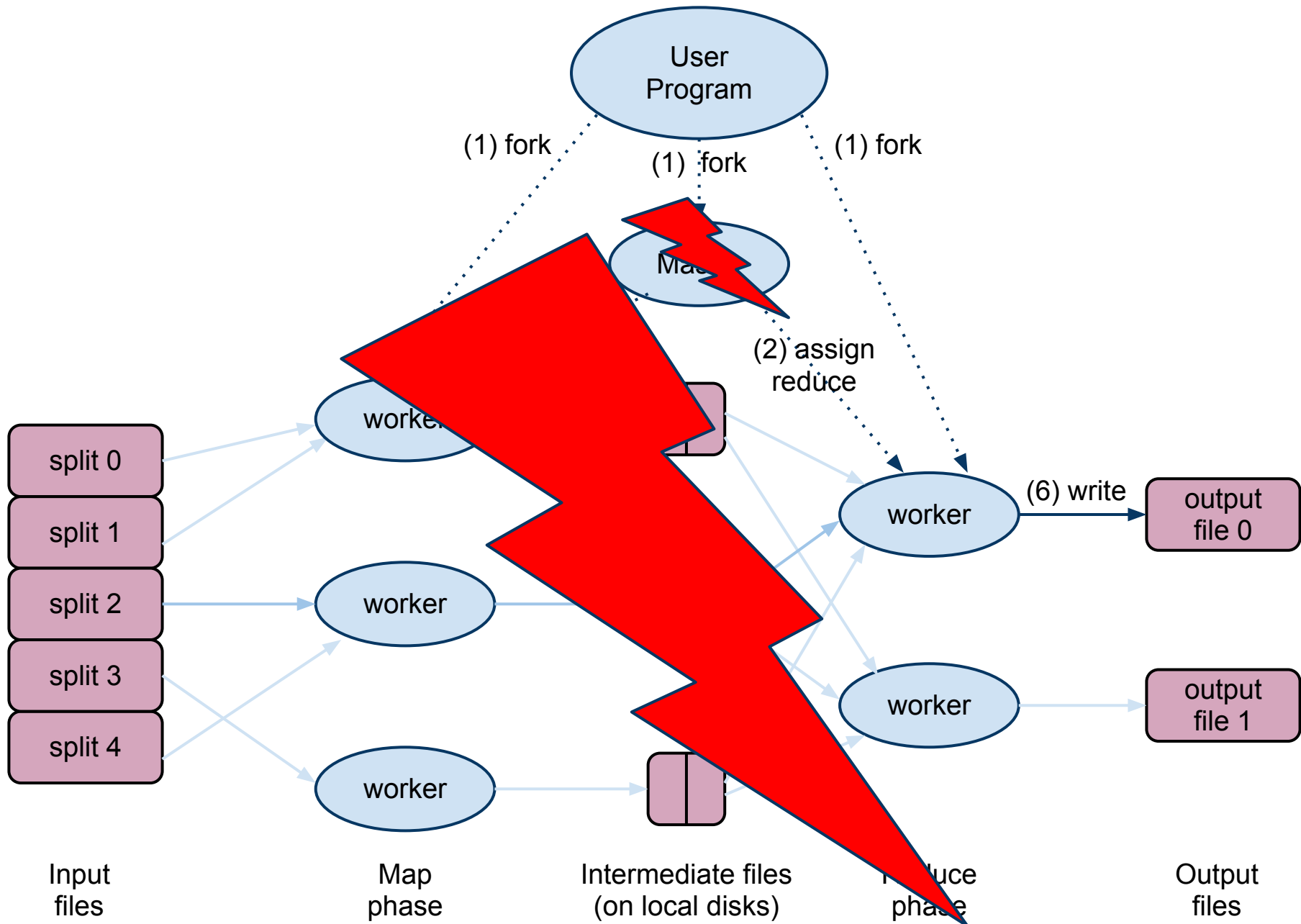
Recover by Checkpointing Map Output



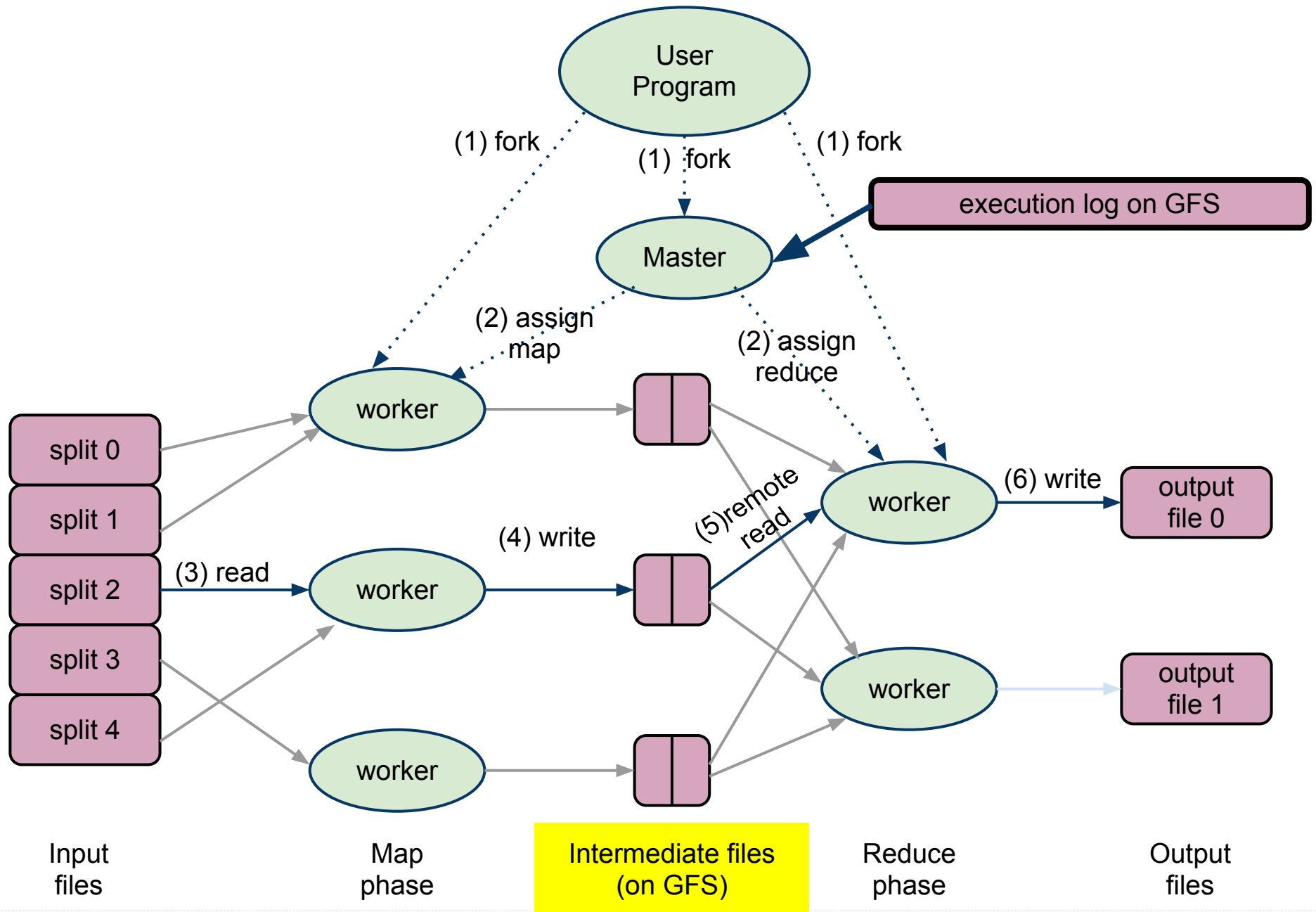
MapReduce: Master Failure



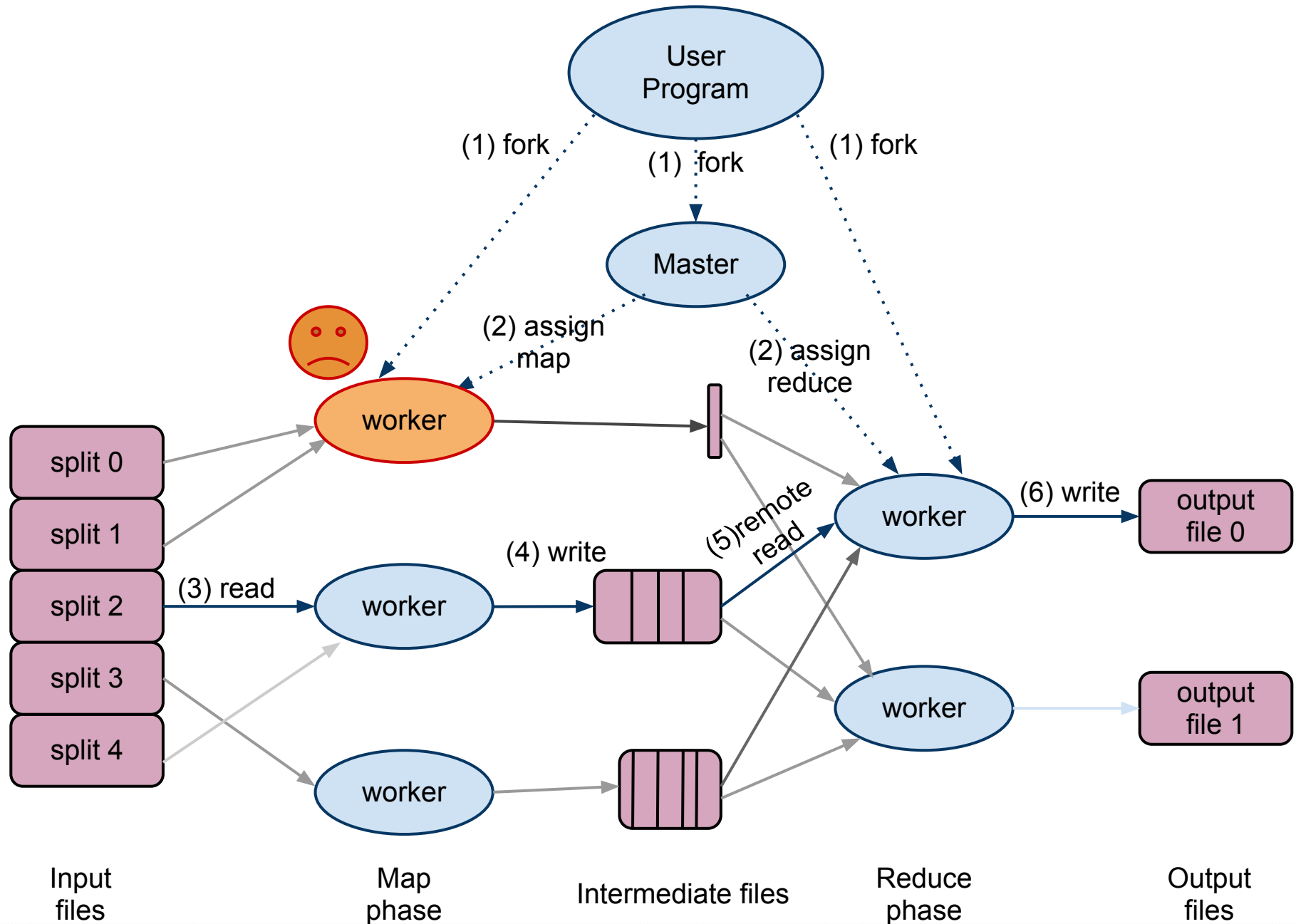
Master as a Single Point of Failure



Resume from Execution Log on GFS



MapReduce: Slow Worker/Task



- Input data is in a partially wrong format or is corrupted
 - Data is mostly well-formatted, but there are instances where your code crashes
 - Corruptions happen rarely, but they are possible at scale
 - Your application depends on an external library which you do not control
 - Which happens to have a bug for a particular, yet very rare, input pattern
 - What would you do?
 - Your job is critical to finish as soon as possible
 - The problematic records are very rare
 - IGNORE IT!
-

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 - Some techniques and tuning tips
 - Usability
-

Performance and Scalability of MapReduce



Terasort and Petasort with MapReduce in Nov 2008

- Not particularly representative for production MRs
- An important benchmark to evaluate the whole stack
- Sorted 1TB (as 10 billion 100-byte uncompressed text) on 1,000 computers in 68 seconds
- Sorted 1PB (10 trillion 100-byte records) on 4,000 computers in 6 hours and 2 minutes

With Open-source Hadoop in May 2009 ([TechReport](#))

- Terasort: 62 seconds on 1460 nodes
 - Petasort: 16 hours and 15 minutes on 3658 nodes
-

Google MapReduce is built upon an set of high performance infrastructure components:

- Google file system (GFS) ([SOSP'03](#))
- Chubby distributed lock service ([OSDI'06](#))
- Bigtable for structured data storage ([OSDI'06](#))
- [Google cluster](#) management system
- Powerful yet energy efficient* hardware and finetuned platform software
- Other house-built libraries and services

Take Advantage of Locality Hints from GFS



- Files in GFS
 - Divided into chunks (default 64MB)
 - Stored with replications, typical $r=3$
 - Reading from local disk is much faster and cheaper than reading from a remote server
 - MapReduce uses the locality hints from GFS
 - Try to assign a task to a machine with a local copy of input
 - Or, less preferable, to a machine where a copy stored on a server on the same network switch
 - Or, assign to any available worker
-

Questions often asked in production:

- How many Map tasks I should split my input into?
- How many Reduce splits I should have?

Implications on scalability

- Master has to make $O(M+R)$ decisions
- System has to keep $O(M*R)$ metadata for distributing map output to reducers

To balance locality, performance and scalability

- By default, each map task is 64MB (== GFS chunksize)
 - Usually, #reduce tasks is a small multiple of #machine
-

- Small map tasks allow fast failure recovery
 - Define "small": input size, output size or processing time
- Big map tasks may force mappers to read from multiple remote chunkservers
- Too many small map shards might lead to excessive overhead in map output distribution

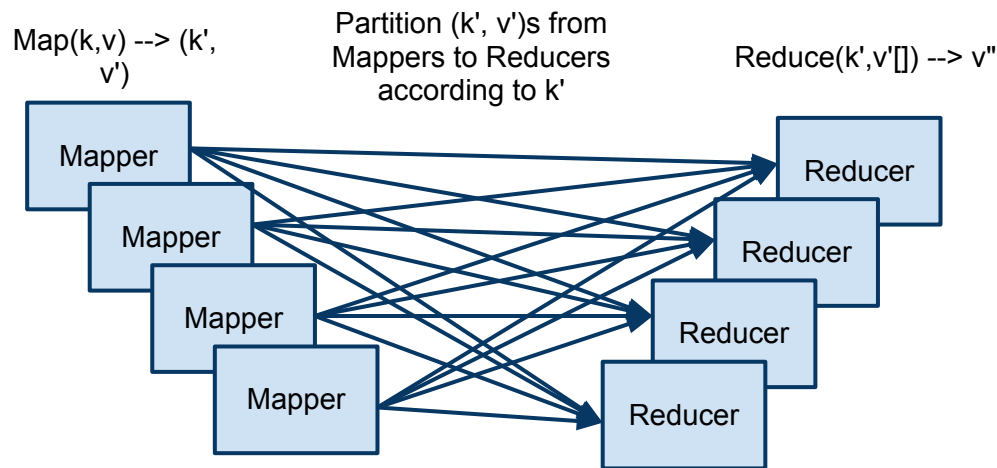
Reduce Task Partitioning Function



It is relatively easy to control Map input granularity

- Each map task is independent

For Reduce tasks, we can tweak the partitioning function instead.



Reduce key	Reduce input size
*.blogspot.com	82.9G
cgi.ebay.com	58.2G
profile.myspace.com	56.3G
yellowpages.superpages.com	49.6G
www.amazon.co.uk	41.7G
average reduce input size for a given key	300K

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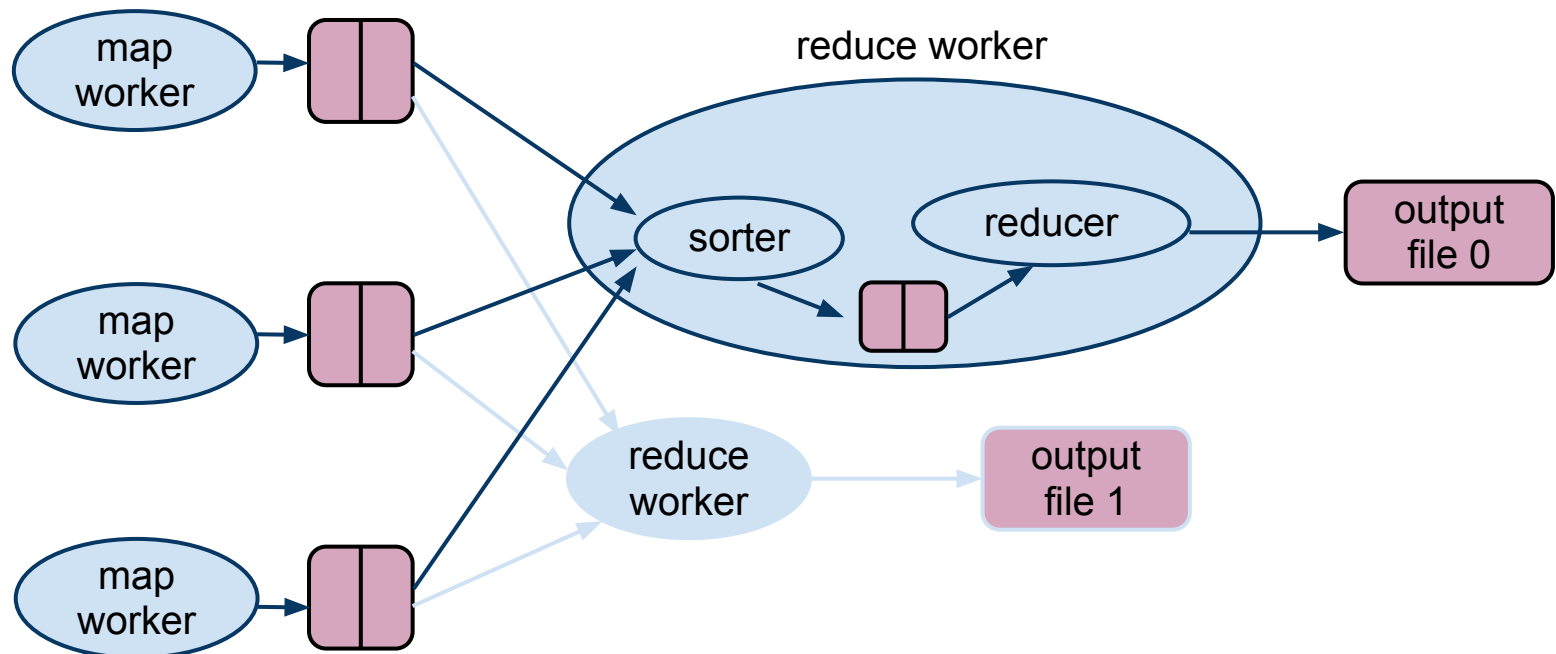
Dealing with Reduce Stragglers



Many reason leads to stragglers but reducing is inherently expensive:

- Reducer retrieves data remotely from many servers
- Sorting is expensive on local resources
- Reducing usually can not start until Mapping is done

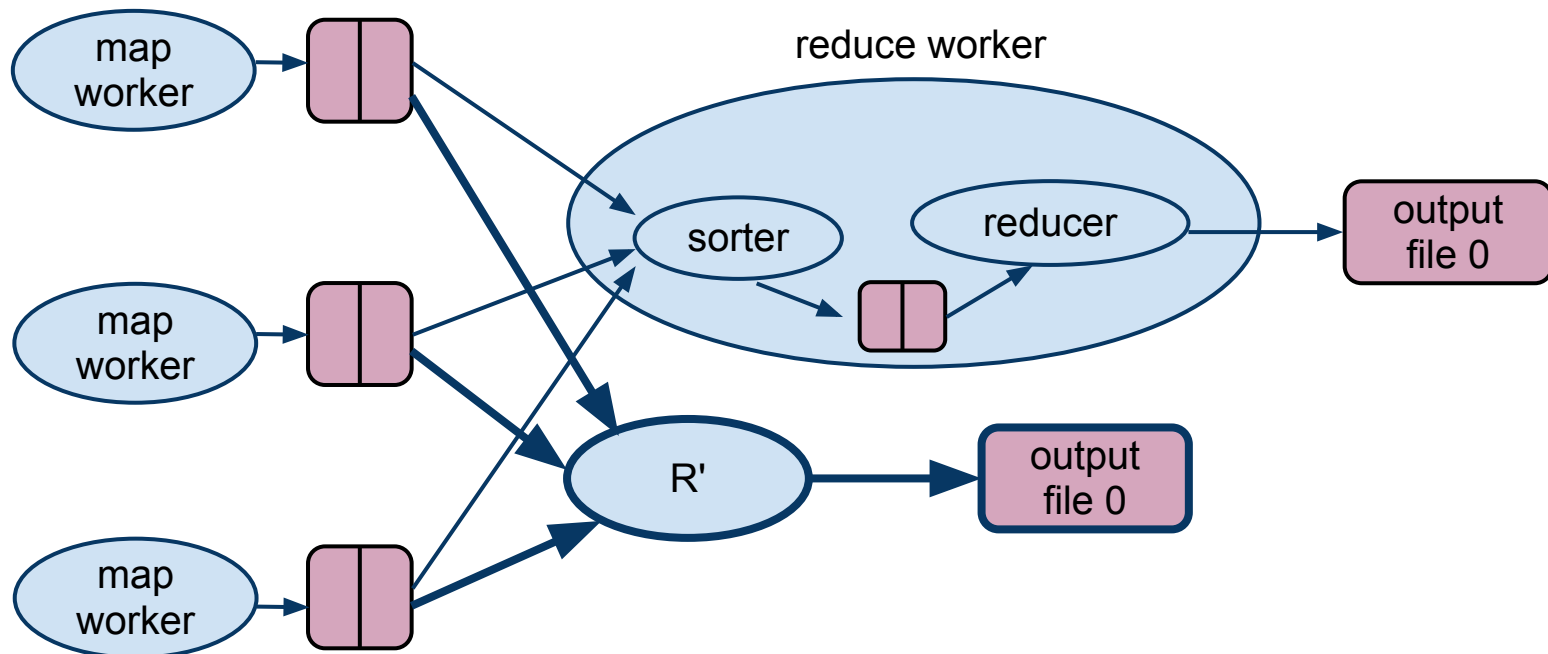
Re-execution due to machine failures could double the runtime.



Dealing with Reduce Stragglers



Technique 1:
Create a backup instance as early and as necessary as possible

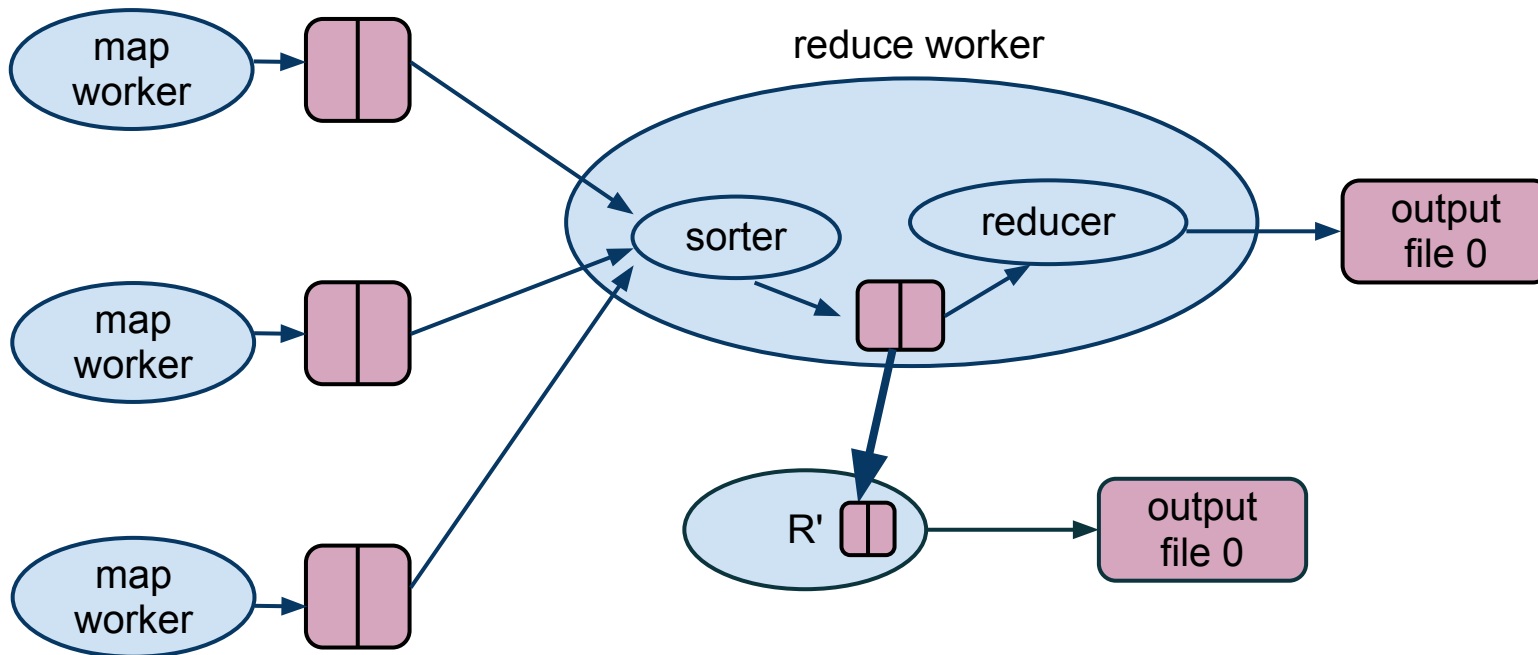


Steal Reduce Input for Backups



Technique 2:

Retrieving map output and sorting are expensive, but we can transport the sorted input to the backup reducer

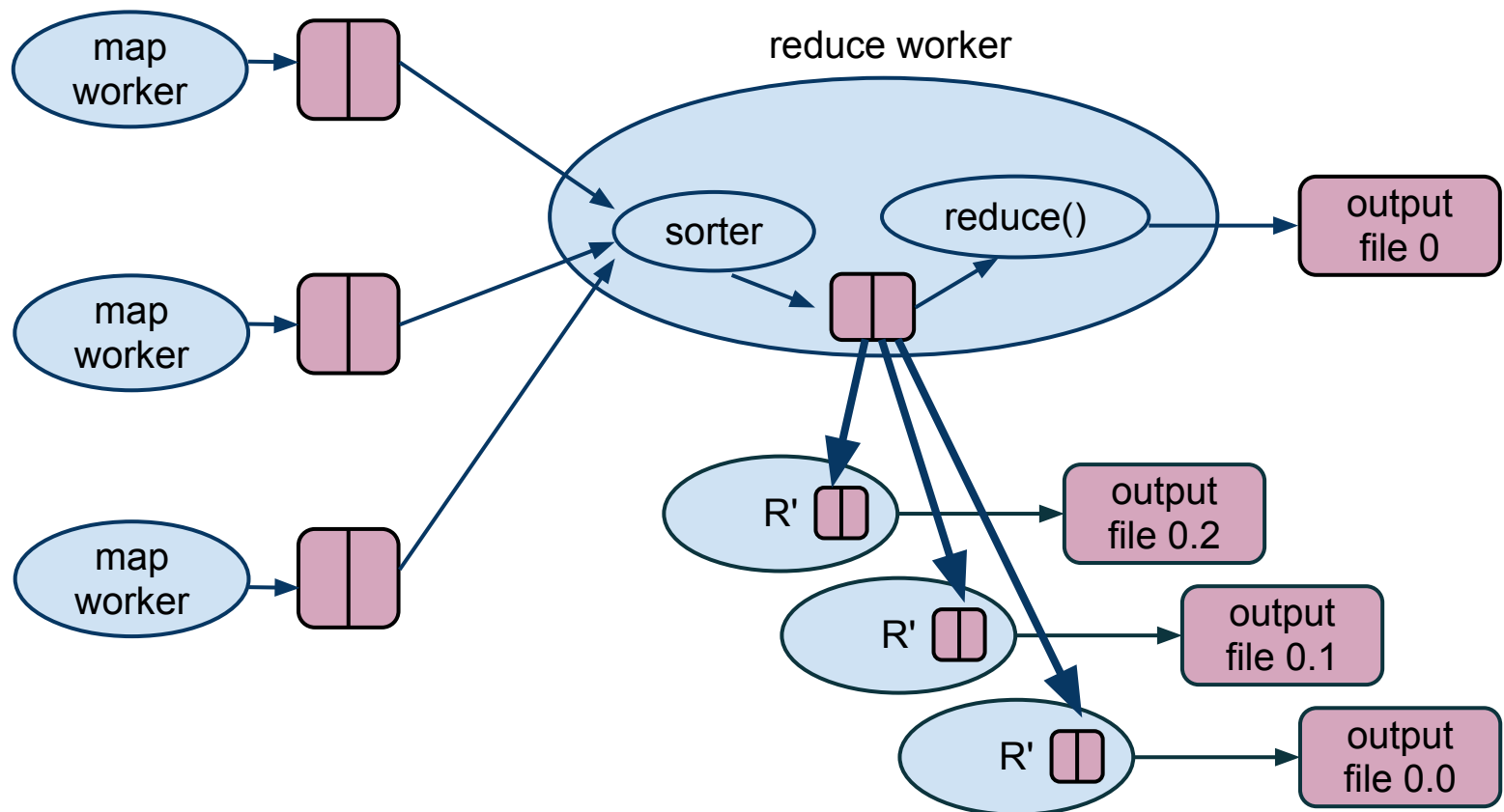


Reduce Task Splitting



Technique 3:

Divide a reduce task into smaller ones to take advantage of more parallelism.



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 - **(Operational) Usability**
 - monitoring, debugging, profiling, etc.
-

Local run mode for debugging/profiling MapReduce applications

Status page to monitor and track progress of MapReduce executions, also

- Email notification
- Replay progress postmortem

Distributed counters used by MapReduce library and application for validation, debugging and tuning

- System invariant
 - Performance profiling
-

Light-weighted stats with only "increment" operations

- per task counters: contributed by each M/R task
 - only counted once even there are backup instances
- per worker counters: contributed by each worker process
 - aggregated contributions from all instances
- Can be easily added by developers

Examples:

- `num_map_output_records == num_reduce_input_records`
 - CPU time spend in `Map()` and `Reduce()` functions
-

Support C++, Java, Python, Sawzall, etc.

Nurtured greatly by Google engineer community

- Friendly internal user discussion groups
- Fix-it! instead of complain-about-it! attitude
- Users contribute to both the core library and contrib
 - Thousands of Mapper Reducer implementations
 - Tens of Input/Output formats
 - Endless new ideas and proposals

- MapReduce is a flexible programming framework for many applications through a couple of restricted Map()/Reduce() constructs
 - Google invented and implemented MapReduce around its infrastructure to allow our engineers scale with the growth of the Internet, and the growth of Google products/services
 - Open source implementations of MapReduce, such as Hadoop are creating a new ecosystem to enable large scale computing over the off-the-shelf clusters
 - MapReduce has many applications for web information retrieval, to parallelize work on large-scale datasets.
-

Thank you!
